

Lecture 14 - dual SVM

Thursday, March 9, 2023 1:32 PM

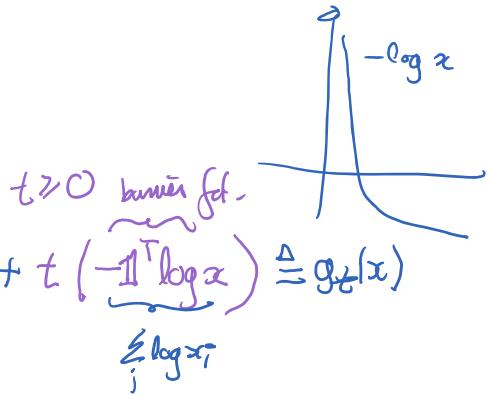
- today :
- IPM
 - duality
 - SVM dual

interior point method :

$$\min f(x)$$

$$\text{s.t. } Ax=b \\ x \geq 0$$

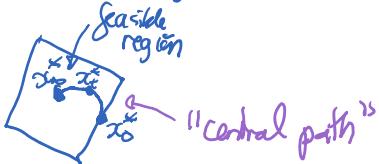
banner method



$$\min f(x) + t \left(-\underbrace{\log x}_{\infty} \right) \triangleq g_t(x)$$

$$\text{s.t. } Ax=b$$

$$\text{dom}(g_t(x)) = \{x : x > 0\}$$



$$x_t^* = \underset{\text{s.t. } Ax=b}{\text{argmin}} g_t(x)$$

Interior point method: Run Newton's method with warm start on $g_t(x)$

and decrease t with a geometric schedule

Convergence: Newton's method takes $\log(\log(\frac{1}{\epsilon})) + \text{cst. } (g_t(x_0) - g_t(x_t^*))$ iterations

to reach ϵ -accuracy

there is no condition if
"self-concordant analysis"

Convergence for IPM: total # of Newton iterations is $\log(\frac{1}{\epsilon}) \cdot \text{some constants}$

↳ dimensionality appears here
but not condition # of f
(vs. gradient methods)

Review of Lagrangian duality

Primal problem

$$p^* \triangleq \inf_x f(x) \text{ s.t. } h_i(x) \leq 0 \quad \forall i \in I$$

Dual problem

$$d^* = \sup_{\alpha \geq 0} d(\alpha)$$

↑ Lagrangian dual set.

$$d(\alpha) \triangleq \inf_x (f(x) + \sum_i \alpha_i h_i(x))$$

Fenchel duality

Fenchel conjugate of f

$$f^*(y) \triangleq \sup_x \langle y, x \rangle - f(x)$$

↑ $x^*(y)$ s.t. $\nabla f(x) = y$

Fenchel duality

$$p^* = \inf_{x \in \mathbb{R}^d} f(x) + g(Ax) \rightarrow \text{e.g. } \sum_{i=1}^m \alpha_i x_i \leq b_i$$

$$\alpha > 0$$

$$d(\alpha) \triangleq \inf_{\mathbf{x}} f(\mathbf{x}) + \sum_i \alpha_i h_i(\mathbf{x})$$

$$P^* = \inf_{\mathbf{x} \in \mathbb{R}^n} f(\mathbf{x}) + g(A\mathbf{x})$$

A: $\mathbb{R}^m \rightarrow \mathbb{R}$
e.g. $\frac{\|x\|_2^2}{2} + \sum_{i=1}^n h_i(x)$

$$d^* = \sup_{y \in \mathbb{R}^m} -f^*(A^T y) - g^*(y)$$

not unique?

saddle pt. perspective of lag. duality

$$\inf_{\mathbf{x}} \underbrace{f(\mathbf{x})}_{\text{s.t. } h_i(\mathbf{x}) \leq 0 \forall i \in I} = \inf_{\mathbf{x}} \left[\sup_{\alpha \geq 0} \underbrace{f(\mathbf{x}) + \sum_i \alpha_i h_i(\mathbf{x})}_{g(\mathbf{x}, \alpha)} \right]$$

Lagrange multiplier

because $g(\mathbf{x}) \triangleq \sup_{\alpha \geq 0} \sum_i \alpha_i h_i(\mathbf{x}) = \begin{cases} +\infty & \text{if } \mathbf{x} \text{ is not feasible} \\ 0 & \text{if } \exists i \text{ s.t. } h_i(\mathbf{x}) > 0 \text{ o.w.} \end{cases}$

$$= S_C(\mathbf{x}) \text{ where } C \triangleq \{ \mathbf{x} : h_i(\mathbf{x}) \leq 0 \forall i \}$$

$$P^* = \inf_{\mathbf{x}} f(\mathbf{x}) + S_C(\mathbf{x})$$

$$= \inf_{\mathbf{x}} \left[\sup_{\alpha \geq 0} g(\mathbf{x}, \alpha) \right]$$

Strong duality $P^* = d^*$

$$\geq \sup_{\alpha \geq 0} \left[\inf_{\mathbf{x}} g(\mathbf{x}, \alpha) \right] \triangleq d^*$$

"weak duality"
not always true

$\triangleq d(\alpha)$
Lagrangian dual fct.

KKT conditions: \Rightarrow necessary conditions for $\mathbf{x}^* \in Q^*$ to be primal & dual optimal
when f & h_i are diff. fcts.

- \mathbf{x}^* is primal feasible (i.e. $h_i(\mathbf{x}^*) \leq 0 \forall i$)
- α^* is dual feasible ($\alpha^* \geq 0$)
- $\nabla f(\mathbf{x}^*) + \sum_i \alpha_i^* \nabla h_i(\mathbf{x}^*) = 0$ [i.e. (\mathbf{x}^*, α^*) stationary pt. of $g(\mathbf{x}, \alpha)$ with respect to \mathbf{x}]
- $\alpha_i^* h_i(\mathbf{x}^*) = 0 \quad \forall i$ "complementary slackness"
 $\hookrightarrow \alpha_i^* > 0 \text{ only on active constraint } h_i(\mathbf{x}^*) = 0$



* if primal problem is convex (i.e. f & h_i 's are convex) and "Slater's condition" \Rightarrow strong duality i.e. $P^* = d^*$

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(i.e. f & h_i are convex) and "Slater's condition" \Rightarrow strong duality i.e. $p^* = d^*$
and KKT conditions are sufficient

$\exists \bar{x}$ rel interior (dom f)
s.t. $h_i(\bar{x}) \leq 0 \forall i$
when h_i is not affine

(e.g. thus supposed \bar{x}^* by solving dual,
then get \bar{x}^* through KKT condition
 $\Rightarrow \bar{x}^*$ solve primal)

↳ see counter-example for strong duality (see Boyd's book example 5.21)

$$\min f(x,y) \\ \text{s.t. } \frac{x^2}{y} \leq 0$$

$$\text{where } f(x,y) = \begin{cases} x^2/y & y > 0 \\ +\infty & \text{o.w.} \end{cases}$$

Violates Slater's condition
and no strong duality

Other comments

- $f(x) \geq d(\alpha) \quad \forall x$ for feasible
- $d(\alpha)$ is always concave in α
- $d(\alpha)$ is obtained through unconstrained opt.
- no unique "Lagrangian dual problem"
e.g. $h_i(\bar{x}) \leq 0 \Leftrightarrow \alpha_i h_i(\bar{x}) \leq 0 \text{ for some } \alpha_i \geq 0$

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SVM struct dual problem

consider primal : $\min_{w, \xi} \frac{\|w\|^2}{2} + \frac{1}{n} \sum_i \xi_i$ $\rightarrow l_i(\tilde{y}) - w^T \phi_i(\tilde{y})$
s.t. $\frac{1}{n} \sum_i \xi_i \geq [H_i(\tilde{y}; w)] \quad \forall \tilde{y} \in \mathcal{Y}, \forall i$

$$L(w, \xi; \alpha) = \frac{\|w\|^2}{2} + \sum_i \xi_i + \sum_{i, \tilde{y}} \alpha_i (\underbrace{l_i(\tilde{y}) - w^T \phi_i(\tilde{y}) - \xi_i}_{h_i(\tilde{y}; w, \xi)})$$

$$\frac{\partial L}{\partial \xi_i} = 0 \Rightarrow \sum_{\tilde{y}} \alpha_i (\phi_i(\tilde{y})) = 1 \quad \text{i.e. } \alpha_i \in \Delta_{|\mathcal{Y}|}$$

$$\nabla_w L = 0 \Rightarrow \lambda w - \sum_{i, \tilde{y}} \alpha_i (\phi_i(\tilde{y})) \nabla \phi_i(\tilde{y}) = 0$$

$$w(\alpha) \triangleq \frac{1}{\lambda n} \sum_{i,j} \alpha_i (\tilde{y}_j) \gamma_i (\tilde{y}_j)$$

KKT condition

$$d(\alpha) = L(w(\alpha), \xi(\alpha); \alpha) = \frac{\lambda \|w(\alpha)\|^2}{2} + \frac{1}{n} \sum_{i,j} \alpha_i (\tilde{y}_j) [l_i(\tilde{y}_j) - w(\alpha)^T \gamma_i (\tilde{y}_j)]$$

notation: $\alpha \in \Delta_{m_1} \times \dots \times \Delta_{m_n}$

$b = (b_i)_{i=1}^n$ $b_i(\tilde{y}) \triangleq \sum_j l_i(\tilde{y}_j)$ $b \in \mathbb{R}^m$ exponentially big
 $w(\alpha) = A\alpha$ where $A \triangleq [\dots \frac{\gamma_1(\tilde{y})}{\lambda n} \dots]$
 $A \in \mathbb{R}^{d \times m}$

$$d(\alpha) = \langle b, \alpha \rangle + \frac{\lambda \|w(\alpha)\|^2}{2} - w(\alpha)^T \left(\sum_{i,j} \alpha_i (\tilde{y}_j) \gamma_i (\tilde{y}_j) \right)$$

$$d(\alpha) = \langle b, \alpha \rangle - \frac{\lambda \|w(\alpha)\|^2}{2}$$

dual problem:

$$\max_{\substack{g_i \in \Delta_{m_i} \\ i \in 1:n}} -\frac{\lambda}{2} \|A\alpha\|^2 + \langle b, \alpha \rangle$$

(exponentially many dual variables)

complementary slackness: $\alpha_i^*(\tilde{y}) > 0 \Rightarrow h_i(\tilde{y}; w^*) = 0$
i.e. $\xi_i^* = l_i(\tilde{y}) - w^*^T \gamma_i(\tilde{y}) = h_i(\tilde{y}; w^*)$

define $\xi_i(\alpha) = \max_{\tilde{y}} h_i(\tilde{y}; w(\alpha))$ by primal feasibility, $\xi_i^* \geq h_i(\tilde{y}; w^*) \forall \tilde{y} \in S_i$
thus $\xi_i^* = \max_{\tilde{y} \in S_i} h_i(\tilde{y}; w^*) = h_i(w^*)$

support vectors are when $\alpha_i^*(\tilde{y}) > 0$ (and $\tilde{y} \neq y^{(i)}$ because $\gamma_i(y^{(i)}) = 0$)
[necessary condition $\tilde{y} \in \arg\max_{\tilde{y}} h_i(\tilde{y}; w^*) \cap S_i$]

II

[necessary condition $\tilde{y} \in \arg\max_{y \in \mathcal{Y}_i} H_i(y; w^*)$]

⇒ sparse solution α^*