

today: properties of FW
AFW alg.

properties of FW

$$1) f(x_t) - \underbrace{\min_{x \in M} f(x)}_{f^*} = O\left(\frac{1}{t}\right)$$

$$2) \text{FW-gap } g_t \geq f(x_t) - f^* \rightarrow \text{certificate of subopt.}$$

$$\min_{t \leq T} g_t \leq O\left(\frac{1}{T}\right) \quad [\text{ie. will stop in } O\left(\frac{1}{\epsilon}\right) \text{ iterations!}]$$

$$3) x_t = p_0^t x_0 + \sum_{u=1}^t p_u^t s_{u-1} \rightsquigarrow x_t \text{ has a } t \text{ "sparse" expansion}$$

in terms of the FW-corners $\sum_{u=0}^{t-1} s_u$

where $\sum_u p_u^t = 1$
 $p_u^t \geq 0$

⊛ "sparse methods"
→ popular in ML

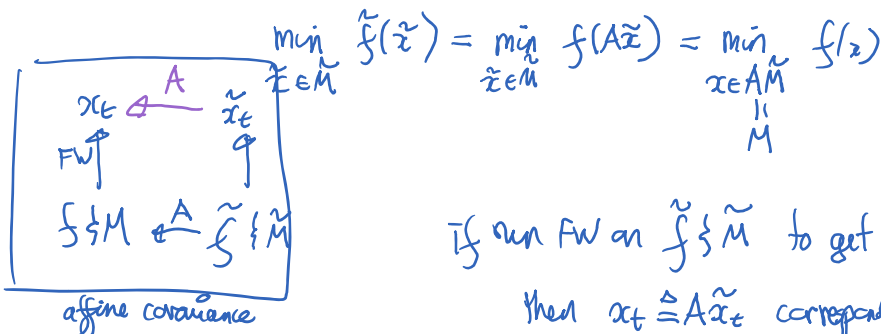
→ see later how to apply on dual of SVM struct

4) there is a $\Omega\left(\frac{1}{t}\right)$ lower bound for FW-like methods for $t \leq d$

5) FW is affine covariant (like Newton's method)

Let $\tilde{M}$ be a new constraint set s.t. $\tilde{M} \xrightarrow{A} M$ linear transformation $\{$ surjective
ie. $M = A\tilde{M}$

define $\tilde{f}(\tilde{x}) \triangleq f(A\tilde{x})$



If run FW on $\tilde{f} \restriction \tilde{M}$ to get $\tilde{x}_t$ as iterates

then $x_t \triangleq A\tilde{x}_t$ corresponds to running FW on $f \restriction M$
(modulo tie breaking)

why? → inner product with gradient

$$\tilde{s}_t = \arg \min_{\tilde{s} \in \tilde{M}} \langle \tilde{s}, \nabla_{\tilde{x}} f(\tilde{x}) \rangle \quad x_t \triangleq A \tilde{x}_t$$

$$\langle \tilde{s}, A^T \nabla_x f(x_t) \rangle$$

$$\langle \underbrace{A \tilde{s}}_{s=A \tilde{s}}, \nabla_x f(x_t) \rangle$$

$$s_t = \arg \min_{s \in M = A \tilde{M}} \langle s, \nabla_x f(x_t) \rangle \Rightarrow \underline{s_t = A \tilde{s}_t} \quad (\text{modulo for breaking})$$

$$\sum_{t=1}^T s_t = A \sum_{t=1}^T \tilde{s}_t \quad \text{set of argmin (ties)}$$

$\Rightarrow$ we want an affine invariant analysis

$$\leadsto C_f \leq L_{\|\cdot\|} (\text{diam}_{\|\cdot\|}(M))^2 \text{ for any } \|\cdot\|$$

affine invariant constant
(we'll see next class)

6) convergence of FW for non-convex

aside: necessary first order condition for constrained optimum

$$\min_{x \in M} f(x)$$

x^* is a local min

$$\Rightarrow \langle \nabla f(x^*), s - x^* \rangle \geq 0 \quad \forall s \in M$$



"stationary pt." for const. opt. problem

$$\Leftrightarrow \min_{s \in M} \langle \nabla f(x^*), s - x^* \rangle \geq 0$$

$$\Leftrightarrow \max_{s \in M} \langle -\nabla f(x^*), s - x^* \rangle \leq 0$$

FW-gap(x^*)

quantify the "non-stationarity"

(if f & M are convex, then it is a sufficient condition for global min)

see Lacoste-Julien 2016 article

$$\min_{s \in M} \text{gap}(x_s) \leq O\left(\frac{1}{\sqrt{t}}\right) \text{ for FW with line search}$$

"non-convex FW"

when f is L -smooth
 M is bounded & convex
but f is not nec. convex

aside: if f is concave, $O\left(\frac{1}{t}\right)$

$\Rightarrow f$ is lower bounded

14h30

away-step FW
comment:

"coordinate"

$$x_t = \sum_u A_u^+ s_u$$

$$x_{t+1} = (1-\gamma_t)x_t + \gamma_t s_t$$

$$= \sum_u \underbrace{A_u^+ (1-\gamma_t)}_{A_u^{t+1}} s_u + \gamma_t s_t$$

$A_u^{t+1} \rightarrow$ previous coordinates shrank by $(1-\gamma_t)$

* FW step moves mass uniformly away from active set $\{s_u\}_{u=1}^t$
to FW corner s_t

$\rightarrow$ unless step-size $\gamma_t = 1$, FW never removes a corner from active set / expansion



$\Rightarrow$ Zig-zag phenomenon close to the boundaries on polytopes

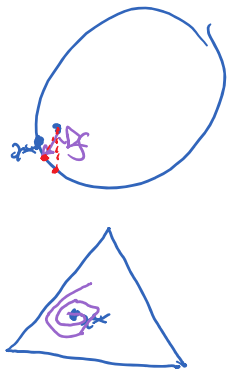
(this is why you get a $\Omega(1/t)$ rate even if f is strongly convex)

$$\left\langle \frac{-\nabla f(x_t)}{\|\nabla f(x_t)\|}, \frac{dx}{dt} \right\rangle \xrightarrow{t \rightarrow \infty} 0$$

* but FW has no problem

an "strongly convex set"
 $\hookrightarrow$ subset set of a strongly convex set.

when set is in the relative interior of M

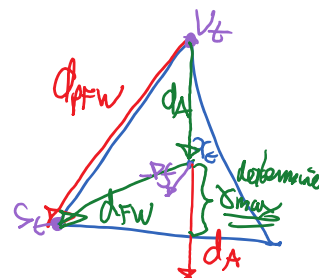


away-step FW fix (solves zig-zagging problem)

in addition to compute FW corner
 $s_t = \arg \min_{s \in M} \langle s, \nabla f(x_t) \rangle$

also compute the "away corner"

$$v_t = \arg \max_{s \in \text{active set}(x_t)} \langle s, \nabla f(x_t) \rangle$$



$\downarrow d_A$

$$d_{FW} = s_t - x_t$$

$$d_A \triangleq x_t - v_t$$

ie. pick d_A if $\underbrace{\langle d_A, -\nabla f(x_c) \rangle}_{\text{"g}_A"} > \underbrace{\langle d_w, -\nabla f(x_c) \rangle}_{\text{"g}_w"}$

$g_A = 0$ when $Df(x)$ is 1
acquire set

where $\gamma_{\max}$ depends on p_i^+ coefficients

$$x_{t+1} = x_t + \gamma_t \left(\frac{d_{x_t}}{x_t - V_t} \right)$$

$$= \sum_u (1 + \gamma_t) \alpha_u s_u - \gamma_t V_t$$

$$(1+\delta)\alpha - \delta \geq 0$$

$$(1 + \gamma_{\max})\alpha - \gamma_{\max} = 0$$

$$(1 + \gamma_{\max})q - \gamma_{\max} = 0 \Rightarrow \boxed{\gamma_{\max} = \frac{q}{1-q}}$$

call this a "drop step" $\rightarrow$ removing V_t from expansion

either you maintain some expansion $\mathcal{A}_t = \sum_i p_i S_i$

you have a feasibility oracle + cutting step oracle
get δ_{max} get V_t

[see Neunig's 2016 paper Moshi's Garden]

1. $\Rightarrow$ (Assume ϵ is $\in (0, \frac{1}{2})$) $\xrightarrow{\text{d}}$ assumption needed for convergence

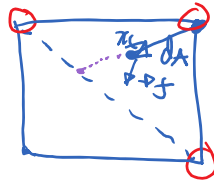
↳ see Nemirovskiy 2016 paper Moshiri & Gordon $\rightarrow$ assumption needed for convergence result

↳ they assume $\text{corner}(M) \subseteq \{0,1\}^d$

and M is described as $Ax=b$

$x \in \{0,1\}^d$

assumption to run alg



AFW has linear convergence rate on polytopes when f is strongly convex
(vs. FW $\rightarrow O(1/\epsilon)$)

③ combine d_{FW} & $d_A \rightarrow$ pairwise FW direction

$$d_{PFW} = d_{FW} + d_A = g_t - x_t + x_t - v_t = g_t - v_t$$

$\underbrace{\hspace{1.5cm}}_{d_{PFW}}$

$$\underbrace{\langle -\nabla f(x_t), d_{PFW} \rangle}_{g_{PFW}} = g_{FW} + g_A \quad \{g_{FW} + g_A \leq 2 \max \{ \dots \} \}$$

$$\text{during AFW: } g_t = \langle -\nabla f(x_t), d_t \rangle = \max \{ g_{FW}, g_A \} \\ \geq \frac{g_{PFW}}{2}$$