

today: graph theory  
DGM

### graphical model

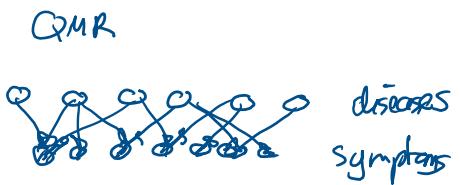
graph. model  $\rightsquigarrow$  prob theory + C.S.  
 R.V.                      graph

graph  $\rightarrow$  efficient data structure

e.g.  $X_1, X_2, \dots, X_n$  R.V.

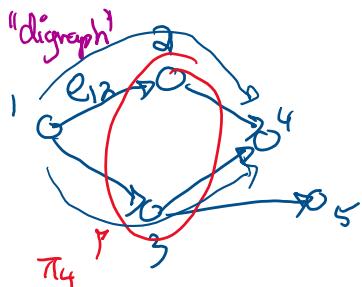
$X_i \in \{0, 1\}$   $n=100s$

$\Rightarrow 2^{100}$  ft's table  
 $\rightsquigarrow$  untractable



### graph theory review:

directed graph  $G = (V, E)$



$$\pi_i^o \triangleq \{j \in V : \exists (j, i) \in E\}$$

set of parents of  $i$

$$\text{children}(i) \triangleq \{j \in V : \exists (i, j) \in E\}$$

$V = \{1, \dots, n\}$  "nodes/vertices"

$E \subseteq V \times V$  "directed edges"

$$e_{12} = (1, 2)$$

directed path  $1 \rightsquigarrow 4$

$$(1, 2), (2, 4)$$

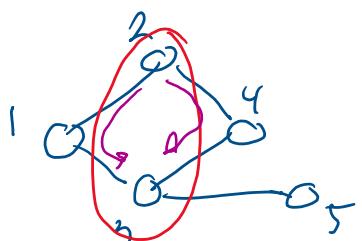
$(1, 3), (3, 4)$

compatible seq.  
of edges

undirected graph :  $G = (V, E)$  where elements of  $E$  are 2-sets  
 (set of 2 different elements of  $V$ )

thus  $\{i, j\} = \{j, i\}$

vs.  $(i, j) \neq (j, i)$  [order matters]



"neighbors" of node 4 or 1

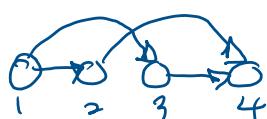
$$N(i) \triangleq \{j \in V : \exists \{i, j\} \in E\}$$

undirected path  $2 \leftarrow 3$

$\oplus$  neighbors replace the parent/children term analogy  
for digraphs

def: DAG = directed acyclic graph  $\triangleq$  digraph with no cycles

def: an ordering  $I: V \rightarrow \{1, 2, \dots, |V|\}$  is said to be topological for digraph G  
 iff nodes in  $\pi_i$  appear before  $i$  in  $I \forall i$



$$(j \in \pi_i \Rightarrow I(j) < I(i))$$

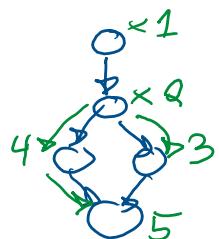
If tp. ordering  $\Rightarrow$  all edges from left to right  
 [no "back edge"]

prop: digraph  $G$  is a DAG  $\Leftrightarrow \exists$  a topological ordering of  $G$

prof:  $\Leftarrow$ ) trivial: no back edge  $\Rightarrow$  no cycle

$\Rightarrow$ ) use DFS algorithm to label nodes in decreasing order  
 when have no children

to construct a tp. sort.  
 in  $O(|E| + |V|)$



tp. sort  $\rightarrow$  finding a topological ordering

### notation for graphical model

n discrete R.V.  $x_1, \dots, x_n$

$\hookrightarrow$  discrete R.V. for simplicity

cond. dist. concept is  
flexible for formulate for cts. R.V.

(see Borel-Kolmogorov paradox)

$V$   $\leftarrow$  set of vertices  
 one R.V. per node

shorthand

joint  $p(x_1=x_1, x_2=x_2, \dots, x_n=x_n) \stackrel{\downarrow}{=} p(x_1, \dots, x_n) \quad x = x_{1:n}$

question: is  $p(x_1, x_2, x_3) \stackrel{?}{=} p(x_2, x_1, x_3) \stackrel{?}{=} p(x_V) \stackrel{?}{=} p(x)$

Yes  $\Rightarrow$  usually in this class  
 we "fixed convention"

for any  $A \subseteq V$

$$p(x_A) = P\{x_i = x_i : i \in A\} = \sum_{x_{A^c}} p(x_A, x_{A^c})$$

subset of "subscript"

$$\text{e.g. } p(x_1, x_2, x_3) = p(x_{\{1, 2, 3\}})$$

Summing over all  
 possible values of  $x_{A^c}$   
 i.e.  $(x_i)_{i \in A^c}$

$$A^c \triangleq V \setminus A$$

### revisit cond. independence

Let  $A, B, C \subseteq V$

$\otimes X_A \perp\!\!\!\perp X_B | X_C$

$$\text{(F) factorization} \Leftrightarrow p(x_A, x_B | x_C) = p(x_A | x_C) p(x_B | x_C) \quad \forall x_A, x_B, x_C \\ \text{s.t. } p(x_C) > 0$$

$$\text{(C) conditioning} \Leftrightarrow p(x_A | x_B, x_C) = p(x_A | x_C) \quad \forall x_A, x_B, x_C \\ \text{s.t. } p(x_B, x_C) > 0$$

$\otimes$  "marginal independence":  $X_A \perp\!\!\!\perp X_B | (\emptyset)$

$$p(x_A, x_B) = p(x_A) p(x_B)$$

$$X \perp\!\!\!\perp Y | Z, W$$

2 facts about cond. indep.

1) can repeat variable in statement  
(for convenience)

$X \perp\!\!\!\perp Y, Z | \underline{Z}, W$  is fine to say  
does not do anything

2) decomposition:  $X \perp\!\!\!\perp (Y, Z) | W \Rightarrow \begin{cases} X \perp\!\!\!\perp Y | W \\ X \perp\!\!\!\perp Z | W \end{cases}$

$\otimes$  pairwise indep.  $\not\Rightarrow$  mutual indep.  $\rightarrow$  see lecture 3  
 $Z = X \oplus Y$   
or

$\otimes$  chain rule  $p(x_V) = \prod_{i=1}^n p(x_i | x_{1:i-1})$

(always true)

last cond.:  $p(x_n | x_{1:n-1})$   
table with  $2^n$  entries

assumption in DBN

$$p(x_V) = \prod_{i=1}^n p(x_i | x_{\pi_i}) \rightarrow \text{tables of } 2^{\max_i |\pi_i| + 1}$$

$$x_i \perp\!\!\!\perp x_{1:i-1} | x_{\pi_i}$$

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directed graph. model

let  $G = (V, E)$  be a DAG

a directed graphical model (DGM) (associated with  $G$ ) (aka a Bayesian network)

is a family of distribution over  $X_V$

$$\mathcal{P}(G) = \left\{ p \text{ is a dist. over } X_V : \exists \text{ legal factors } f_i \text{'s s.t. } p(x_V) = \prod_{i=1}^n f_i(x_i | x_{\pi_i}) \quad \forall x \right\}$$

$p(x_v) = \prod_{i=1}^n f_i(x_i | x_{\pi_i})$

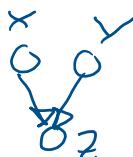
st.  $f_i : \mathcal{L}_{x_i} \times \mathcal{L}_{x_{\pi_i}} \rightarrow [0, 1]$

st.  $\sum_{x_i} f_i(x_i | x_{\pi_i}) = 1 \quad \forall x_{\pi_i}$

$\Rightarrow f_i$  is like a CPT  
(conditional prob. table)

Terminology: if we can write  $p(x_v) = \prod_i f_i(x_i | x_{\pi_i})$  this defines  $G$   
then we say that " $p$  factorizes according to  $G$ " &  $p \in \mathcal{F}(G)$

one example



$$p \in \mathcal{F}(G) \Leftrightarrow \exists f_x, f_y, f_z \text{ legal}$$

$$\text{st. } p(x_1, x_2, x_3) = f_x(x_1) f_y(x_2 | x_1) f_z(x_3 | x_2)$$

"leaf plucking" property (fundamental property of DGM)

Let  $p \in \mathcal{F}(G)$ ; let  $n$  be a leaf in  $G$  (i.e.  $n$  has no children)

a) then  $p[x_{1:n-1}] \in \mathcal{F}(G - \{n\})$

means remove node  $n$  from  $G$

b) moreover, if  $p(x_{1:n}) = \prod_{i=1}^n f_i(x_i | x_{\pi_i})$  then  $p(x_{1:n-1}) = \prod_{i=1}^{n-1} f_i(x_i | x_{\pi_i})$

Proof:  $p(x_n, x_{1:n-1}) = f_n(x_n | x_{\pi_n}) \prod_{i=1}^{n-1} f_i(x_i | x_{\pi_i})$  no  $n$  in any  $\pi_i$

$$p(x_{1:n-1}) = \sum_{x_n} p(x_n, x_{1:n-1}) = \left[ \sum_{x_n} f_n(x_n | x_{\pi_n}) \right] \left( \prod_{i=1}^{n-1} f_i(x_i | x_{\pi_i}) \right)$$

1 by defin

proposition: if  $p \in \mathcal{F}(G)$

let  $\{f_i\}$  be a factorization w.r.t.  $G$  of  $p$

then  $\forall i$ ,  $p(x_i | x_{\pi_i}) = f_i(x_i | x_{\pi_i}) \quad \forall x_{\pi_i}$

i.e. factors are correct conditionals of  $p$   
(and are unique)

prof: WLOG, let  $(1, \dots, n)$  be a top. sorting of  $G$  (possible since  $G$  is a DAG)

- let  $i$  be fixed; we want to show that  $p(x_i | x_{\pi_i}) = f_i(x_i | x_{\pi_i})$

- $n$  is a leaf  $\rightarrow$  "pluck it" to get  $p(x_{1:n-1}) = \prod_{i=1}^{n-1} f_i(x_i | x_{\pi_i})$

- well we justed, we want to show that  $p(x_{1:n}) = \prod_{i=1}^n p(x_i | x_{1:i-1})$
- $n$  is a leaf  $\rightarrow$  "pluck it" to get  $p(x_{1:n-1}) = \underbrace{\prod_{j=1}^{n-1} f_j}_{\text{factorization for a DAG where } x_{n-1} \text{ is}} \text{ bsg } | x_{n-1}$

• pluck  $n-1$  as well? (as leaf, because of top. sort.)

• keep doing this [in reverse order  $n-1, n-2, \dots, 1+1$ ] to get

$$p(x_{1:i}) = f_i(x_i | x_{\pi_i}) \prod_{j < i} f_j(x_j | x_{\pi_j})$$

$$\stackrel{\Delta}{=} g(x_{1:i-1}) \quad \begin{matrix} \text{i.e. there is no } x_i \\ \text{by top. sort. property} \end{matrix}$$

partition  $1:i$  as  $\{x_i\} \cup \pi_i \cup A$  [ $A \triangleq 1:i-1 \setminus \pi_i$ ]

$$p(x_{1:i}) = p(x_i, x_{\pi_i}, x_A)$$

$$p(x_i | x_{\pi_i}) = \frac{\sum_{x_A} p(x_i, x_{\pi_i}, x_A)}{\sum_{x_i} \sum_{x_A} p(x_i, x_{\pi_i}, x_A)} = \frac{f_i(x_i | x_{\pi_i}) \sum_{x_A} g(x_{1:i-1})}{\sum_{x_i} \left( \sum_{x_A} f_i(x_i | x_{\pi_i}) g(x_{1:i-1}) \right)} = f_i(x_i | x_{\pi_i})$$

(now), we can safely write

$$\mathcal{G}(G) = \{p : p(x) = \prod_{i=1}^n p(x_i | x_{\pi_i})\}$$

properties of DGM:

adding edges  $\Rightarrow$  more distributions

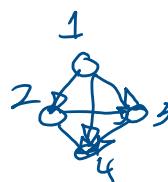
i.e.  $G = (V, E)$  and  $G' = (V, E')$  with  $E' \supseteq E$   
then  $\mathcal{G}(G) \subseteq \mathcal{G}(G')$

examples:

•  $E = \emptyset$  (trivial graph)  $\Rightarrow \mathcal{G}(G)$  contains only fully independent distributions  
i.e.  $p(x_V) = \prod_{i=1}^n p(x_i) \not\in \mathcal{G}(G)$

• complete DAG

forall  $i \neq j$  either  $(i, j) \in E$   
 $\text{or } (j, i) \in E$



get wlog

$$p(x_V) = \prod_{i=1}^n p(x_i | x_{1:i-1})$$

Draw 1 if either  $(i,j) \in E$   
or  
 $(j,i) \in E$



$$p(x_V) = \prod_{i=1}^n p(x_i | x_{1:i-1})$$

(like chain rule)

$\Rightarrow$  all dist. on  $x_V$   
are in  $\mathcal{G}$  (complete)

⊕ removing edges impose more factorization constraints

(and thus more cond. indeps. assumptions)

"absence of edge" is meaningful in a DGM