

today:

- max-product alg.
- junction tree
- HMM

getting conditions

$x_E = \bar{x}_E$
 $\downarrow$
 $p(x_i | x_E) \propto p(x_i, \bar{x}_E)$
 indicates values we are conditioning on
 keep this fixed during marginalization for each $j \in E$

(formal trick): redefine $\tilde{\psi}_j(x_j) \triangleq \psi_j(x_j) \cdot \delta(x_j, \bar{x}_j)$
 Kronecker delta fct.
 $\delta(a, b) = \begin{cases} 1 & \text{if } a=b \\ 0 & \text{o.w.} \end{cases}$

$$\text{computing } m_{j \rightarrow i}(x_i) = \sum_{x_j} \tilde{\psi}_j(x_j) \text{stuff}(x_i, x_j) \\ \rightarrow \psi_j(\bar{x}_j) \text{stuff}(x_i, \bar{x}_j)$$

at the end, result of sum-product

$$\text{will give } p(x_i, \bar{x}_E) = \frac{1}{Z_i} \psi_i(x_i) \prod_{k \in N(i)} m_{k \rightarrow i}(x_i)$$

renormalize over x_i to get $p(x_i | \bar{x}_E)$

max-product alg.

for sum-product, main property used was distributivity of $\oplus$ over $\odot$

all need is that $(R, \oplus, \odot)$ is $[a \odot b \oplus a \odot c = a \odot (b \oplus c)]$

semi-ring $\rightarrow$ ring that doesn't require additive inverses

⊛ can do "sum-product" on other semi-rings

"in this case, ..."

⊛ can do "sum-product" on other semi-rings

$$(\mathbb{R} \cup \{-\infty\}, \max, +) \quad \max(a+b, a+c) = a + \max(b, c)$$

$$(\mathbb{R}_+, \max, \cdot) \quad \max(a \cdot c, b \cdot c) = a \cdot \max(b, c)$$

→ "max-product"

$$\max_{x_{1:n}} \prod_i f_i(x_i) = \prod_{i=1}^n \max_{x_i} f_i(x_i)$$

$$m_{i \rightarrow j}(x_j) = \max_{x_i} \left[\psi_i(x_i) \psi_{ij}(x_i, x_j) \prod_{k \in N(i) \setminus \{j\}} m_{k \rightarrow i}(x_i) \right]$$

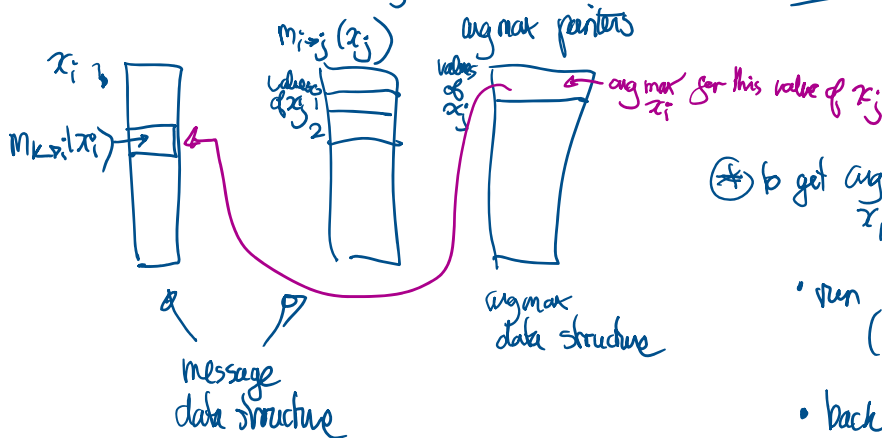
⚡
x_i

→ for getting argmax

store argument of this max as a fct. of x_j



$$\max_{x_{1:5}} \prod_c \psi_c(x_c) = \frac{1}{Z} \max_{x_5} [m_{4 \rightarrow 5}(x_5) \psi_5(x_5)]$$



⊛ to get argmax $p(x_{1:n})$ "decoding"

- run max-product alg. (only forward messages)
- backtrack the argmax pointers to get full argmax

aka. Viterbi alg.

property of tree UGM
 $p \in \mathcal{G}(\text{tree})$
 with non-zero marginals

$$\Rightarrow p(x) = \prod_{i \in V} \psi_i(x_i) \prod_{(i,j) \in E} \frac{\psi_{ij}(x_i, x_j)}{p(x_i)p(x_j)}$$

↓

$$\frac{p(x_i|x_j)}{p(x_i)}$$

⊗ Similar to DBM; for any set of factors $\{f_{ij}(x_i, x_j)\}$
 $\{f_i(x_i)\}$ $f_{ij} \geq 0$
 $f_i \geq 0$

st. 'local consistency property'

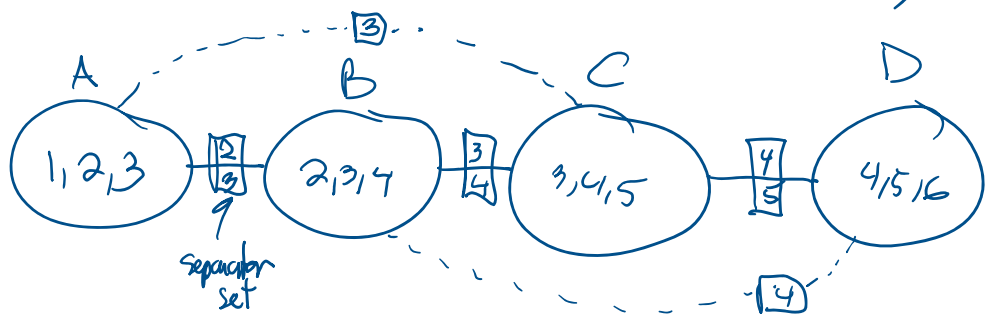
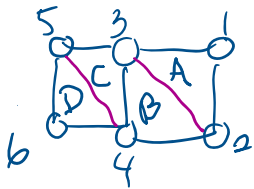
$$\begin{cases} \sum_{x_i} f_{ij}(x_i, x_j) = f_j(x_j) & \forall x_j \\ \sum_{x_j} f_{ij}(x_i, x_j) = f_i(x_i) & \forall x_i \end{cases}$$

$$\sum_{x_i} f_i(x_i) = 1$$

then if define joint $p(x) = \frac{1}{Z} \prod_i f_i(x_i) \prod_{f_{ij} \in E} \frac{f_{ij}(x_i, x_j)}{f_i(x_i) f_j(x_j)}$

(for tree) $\rightarrow$ then can show that $Z=1$
 and f_{ij} are correct marginals
 eg $p(x_i) = f_i(x_i)$
 etc.. //

junction tree alg: generalization of sum-product to a clique tree (with J.T. property)



above is a clique tree with the "running intersection property" $\leftarrow$ "junction tree"

$\hookrightarrow$ if $j \in C_1 \cap C_2$, then $j \in C_k \forall C_k$ along path from C_1 to C_2

to build a J.T. on a Δ -graph:
 • use maximum weight spanning tree alg on clique graph (with size of separator set as weight on edges)

$\Rightarrow$ has running intersection property

$\exists$ a J.T. $\Leftrightarrow$ triangulated / decomposable graph

$\Leftarrow$ running graph

$\exists$ a J.T. $\Leftrightarrow$ triangulated / decomposable graph

4 $\rightarrow$ running graph eliminate

⊗ when have J.T., one can show

$$p(x_v) = \frac{\prod_{c \in C} \psi_c(x_c)}{\prod_{s \in S} \psi_s(x_s)}$$

separator sets of J.T.

$\forall \{1, 2, 3\}$
 $C_1 \{1, 2\}$
 $C_2 \{2, 3\}$

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J.T. alg: reconstruct the joint formulation

by starting with $p(x_v) = \frac{1}{Z} \frac{\prod_c \psi_c(x_c)}{\prod_s \psi_s(x_s)}$

where $\psi_s(x_s) = 1$ at initialization

do message passing on J.T. to update

ψ_c^{new}
 ψ_s^{new}

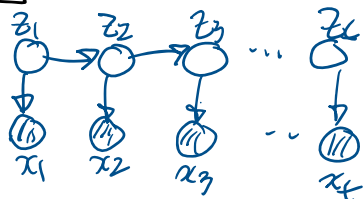
at end of Message passing

$\psi_c^{\text{new}} \rightarrow p(x_c)$
 $\psi_s \sim p(x_s)$

• which leaves the joint invariant

(see Mike's book ch. 17 for details)

HMM (Hidden Markov model)

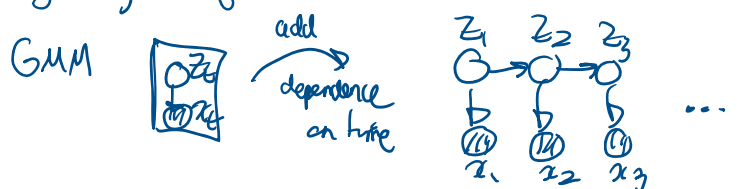


$z_t \in \{1, \dots, K\}$ discrete

(data $z_t \sim$ Gaussian) $\rightarrow$ Kalman filter

$x_t \leftarrow$ obs. eg. speech sound waves
discrete eg. DNA sequence

HMM $\rightarrow$ generalization of mixture model



$$\text{DGM: } p(x_{1:T}, z_{1:T}) = p(z_1) \prod_{t=1}^T p(x_t | z_t) \prod_{t=1}^T p(z_t | z_{t-1})$$

$$\text{DGM: } p(x_{1:T}, z_{1:T}) = p(z_1) \underbrace{\prod_{t=1}^T p(x_t | z_t)}_{\text{emission prob.}} \underbrace{\prod_{t=2}^T p(z_t | z_{t-1})}_{\text{transition prob.}}$$

often: emission & trans. prob. are homogeneous in time
(i.e. do not depend on t)

$$p_t(x_t | z_t) = f(x_t | z_t)$$

$$p_t(z_t = i | z_{t-1} = j) = A_{ij}$$

"stochastic matrix"

$$A = \begin{pmatrix} \vdots & \vdots & \vdots \\ \vdots & \vdots & \vdots \\ \vdots & \vdots & \vdots \end{pmatrix} \text{ def. over } z_t$$

$$\sum_i A_{ij} = 1$$

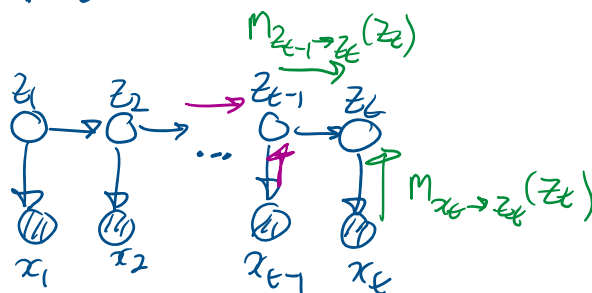
inference tasks:

- prediction $p(z_t | x_{1:t-1})$ "where next"
- filtering $p(z_t | x_{1:t})$ "where now"
- smoothing $p(z_t | x_{1:T})$ "where in the past"

$T \geq t$

α -recursion

let's run sum-product alg. to derive recursion to compute $p(z_t | x_{1:t})$



to compute $p(z_t | \bar{x}_{1:t})$ or $p(z_t | \bar{x}_{1:t})$

$$\underbrace{p(z_t | \bar{x}_{1:t})}_{= \alpha_t(z_t)} = \frac{1}{Z} \sum_z 1 \underbrace{m_{z_{t-1} \rightarrow z_t}(z_t)}_{m_{z_{t-1} \rightarrow z_t}(z_t)} \underbrace{m_{x_t \rightarrow z_t}(z_t)}_{m_{x_t \rightarrow z_t}(z_t)} \quad (\text{here } z=1)$$

$$m_{x_t \rightarrow z_t}(z_t) = \sum_{x_t} p(x_t | z_t) \underbrace{\delta(x_t, \bar{x}_t)}_{\text{conditioning free}} = p(\bar{x}_t | z_t)$$

$$m_{z_{t-1} \rightarrow z_t}(z_t) = \sum_{z_{t-1}} p(z_t | z_{t-1}) \underbrace{m_{z_{t-2} \rightarrow z_{t-1}}(z_{t-1})}_{m_{z_{t-2} \rightarrow z_{t-1}}(z_{t-1})} m_{x_{t-1} \rightarrow z_{t-1}}(z_{t-1})$$

$$p(z_t | \bar{x}_{1:t})$$

$$\alpha_t(z_t) = p(\bar{x}_t | z_t) \sum_{z_{t-1}} p(z_t | z_{t-1}) \alpha_{t-1}(z_{t-1})$$

$$\underbrace{p(z_{t-1} | \bar{x}_{1:t-1})}_{\alpha_{t-1}(z_{t-1})}$$

$$\boxed{\alpha_t(z_t) = \underbrace{p(\bar{x}_t | z_t)}_{\text{vector}} \sum_{z_{t-1}} \underbrace{p(z_t | z_{t-1})}_{\text{matrix}} \underbrace{\alpha_{t-1}(z_{t-1})}_{\text{vector}}} \quad \alpha_{t-1}(z_{t-1})$$

α -recursion aka "forward recursion"

like the "collect phase" in sum-product alg.

initialization: $\alpha_1(z_1) = p(z_1, \bar{x}_1) = p(\bar{x}_1 | z_1) p(z_1)$

Let $\alpha_t(z_t) \triangleq p(\bar{x}_t | z_t)$

$$\boxed{\alpha_t = \alpha_t \odot (A \alpha_{t-1})}$$

↑
Hadamard product

$$\tilde{\alpha}_t(z_t) \triangleq p(z_t | \bar{x}_{1:t}) \quad \text{"filling distribution"}$$

space complexity: $O(k)$ extra storage

time complexity: $O(tk^2)$

$$\begin{aligned} &= \frac{\alpha_t(z_t)}{\sum_{z_t} \alpha_t(z_t)} \} \sum_{z_t} p(z_t, \bar{x}_{1:t}) \\ &= p(\bar{x}_{1:t}) \\ &\quad \text{"evidence prob."} \end{aligned}$$