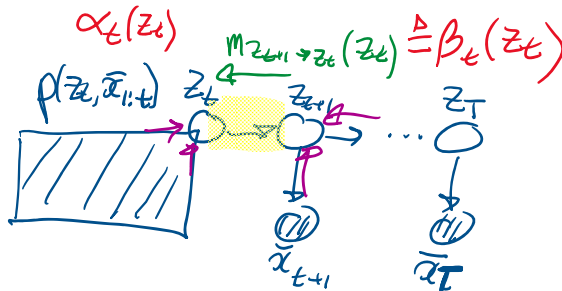


today: EM for HMM
· info theory & KL

β -recursion: (smoothing)



$$p(z_t, \bar{x}_{1:T}) = \frac{1}{Z} \alpha_t(z_t) \underbrace{m_{z_{t+1} \rightarrow z_t}(z_t)}_{\triangleq \beta_t(z_t)}$$

$$\underbrace{m_{z_{t+1} \rightarrow z_t}(z_t)}_{\beta_t(z_t)} = \sum_{z_{t+1}} p(z_{t+1} | z_t) p(\bar{x}_{t+1} | z_{t+1}) \underbrace{m_{z_{t+2} \rightarrow z_{t+1}}(z_{t+1})}_{\beta_{t+1}(z_{t+1})}$$

$$\boxed{\beta_t(z_t) = \sum_{z_{t+1}} p(z_{t+1} | z_t) p(\bar{x}_{t+1} | z_{t+1}) \beta_{t+1}(z_{t+1})}$$

β -recursion (aka. backward recursion)

turns out that $\boxed{\beta_t(z_t) \triangleq p(\bar{x}_{t+1:T} | z_t)}$

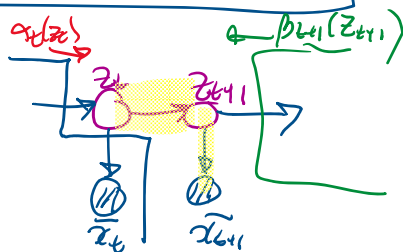
why? $p(z_t, \bar{x}_{1:T}) = p(\bar{x}_{1:T} | z_t) p(z_t)$

$$\stackrel{\text{c.i.d.}}{=} p(\bar{x}_{t+1:T} | z_t) \underbrace{p(\bar{x}_{1:t} | z_t) p(z_t)}_{\alpha_t(z_t)}$$

$\Rightarrow \beta_t(z_t)$

initialization $\beta_T(z_T) = 1 \forall z_T$

edge marginal:



$$\boxed{p(z_t, z_{t+1}, \bar{x}_{1:T}) = \alpha_t(z_t) \beta_{t+1}(z_{t+1}) p(z_{t+1} | z_t) p(\bar{x}_{t+1} | z_{t+1})}$$

numerical stability trick:

numerical stability trick:

issue: α_t & β_t can easily go to $1e-1000$

two possibilities:

a) (general) store $\log(\alpha_t)$ instead

$$\begin{aligned}\log\left(\sum_i a_i\right) &= \log\left(\tilde{\alpha} \left(\sum_i \frac{a_i}{\tilde{\alpha}}\right)\right) \quad (a_i > 0) \\ \text{call } \tilde{\alpha} &\triangleq \max_i a_i &= \log \tilde{\alpha} + \log\left(\sum_i \frac{a_i}{\tilde{\alpha}}\right) \\ i_{\max} &\triangleq \arg\max_i a_i &\log\left(\sum_i \exp(\log a_i - \log(\tilde{\alpha}))\right) \\ \log(\tilde{\alpha}) &= \log(a_{i_{\max}}) = \max_i \log(a_i)\end{aligned}$$

$$\log\left(\sum_i a_i\right) = \log(\tilde{\alpha}) + \log\left(1 + \sum_{i \neq i_{\max}} \exp(\log(a_i) - \log(\tilde{\alpha}))\right)$$

b) normalize the message

• α -recursion, use $\tilde{\alpha}_t(z_t) \triangleq p(z_t | \bar{x}_{1:t})$

$$\text{before } \alpha_t = \alpha_t \odot A \alpha_{t-1} \quad \tilde{\alpha}_t = \frac{\alpha_t \odot A \alpha_{t-1}}{\sum_{z_t} (\quad)} \triangleq c_t$$

$$\begin{aligned}\text{you can show that } c_t &= \sum_{z_t} (\alpha_t \odot A \alpha_{t-1})(z_t) \\ &= p(\bar{x}_t | \bar{x}_{1:t-1})\end{aligned}$$

$$p(\bar{x}_{1:T}) = \prod_{t=1}^T p(\bar{x}_t | \bar{x}_{1:t-1}) = \prod_{t=1}^T c_t$$

• β -recursion:

$$\text{define } \tilde{\beta}_t(z_t) = \frac{p(\bar{x}_{t+1:T} | z_t)}{p(\bar{x}_{t+1:T} | \bar{x}_{1:t})} \quad \prod_{t=t_0+1}^T c_t$$

note: $\sum_{z_t} \tilde{\beta}_t(z_t) \neq 1$

exercise: derive $\tilde{\beta}$ -recursion

15h21

ML for HMM

for some parametric model for dist. on z_t

ML for HMM

for some parametric model for dist. on x_t

• suppose $p(x_t | z_t = k) = f(x_t | \eta_k)$ e.g. Gaussian on x_t

$$\eta = (\eta_k)_{k=1}^K$$

• $p(z_{t+1} = i | z_t = j) = A_{ij}$

• $p(z_1 = i) = \pi_i$

$$\Theta = (\eta, A, \pi)$$

want to estimate $\hat{\eta}, \hat{A}, \hat{\pi}$ by ML from data $x = (x^{(i)})_{i=1}^N$

$$x^{(i)} = x_{1:T}^{(i)}$$

→ use EM at s^{th} iteration

E step: $q_{s+1}(z) = p(z | x, \Theta^{[s]})$

M step: $\Theta^{[s+1]} = \underset{\Theta \in \Theta}{\text{argmax}} \mathbb{E}_{q_{s+1}(z)} [\log p(x, z | \Theta)]$

complete log-likelihood

$$\log p(x, z | \Theta) = \sum_{i=1}^N \left(\underbrace{(\log p(z^{(i)}))}_{\sum_k z_{1,k}^{(i)} \log \pi_k} + \sum_{t=1}^{T_i} \underbrace{\log p(\bar{x}_t^{(i)} | z_t^{(i)})}_{\sum_k z_{t,k}^{(i)} \log f(\bar{x}_t^{(i)} | \eta_k)} + \sum_{t=2}^{T_i} \underbrace{\log p(z_t^{(i)} | z_{t-1}^{(i)})}_{\sum_{l,m} z_{t,l}^{(i)} z_{t-1,m}^{(i)} \log A_{lm}} \right)$$

↑
huge variable

$\mathbb{E}_{q_{s+1}} [\log p(x, z | \Theta)]$

$\mathbb{E}_{q_{s+1}} [z_{t,k}^{(i)}] = q_{s+1}(z_{t,k}^{(i)} = 1) \triangleq \gamma_{t,k}^{(i)}$ ← obtained using O-P recursion on i^{th} time series

smoothing def: $p(z_{t,k}^{(i)} = 1 | x_{1:T}^{(i)}; \Theta^{[s]})$

$q_{s+1}(\bar{z}_{t,l}^{(i)} = 1, z_{t-1,m} = 1) = p(z_{t,l}^{(i)} = 1, z_{t-1,m} = 1 | \bar{x}_{1:T_i}^{(i)}; \Theta^{[s]})$

$\triangleq \gamma_{t,l,m}^{(i)}$

A matrix $\begin{pmatrix} m \\ l \end{pmatrix}$

note: you should have

$$\sum_l \gamma_{t,l,m}^{(i)} = \gamma_{t-1,m}^{(i)}$$

smoothing edge marginal probs in HMM (O-P recursion)

max. with respect to Θ :

$$\hat{\pi}_k^{[s+1]} = \sum_{i=1}^N \sum_{z_1} \gamma_{1,k}^{(i)}$$

$$\hat{A}_{lm}^{[s+1]} = \sum_{i=1}^N \sum_{t=2}^{T_i} \gamma_{t,l,m}^{(i)}$$

$\hat{\eta}_k \rightarrow$ soft count ML

$$\tau_k^{(i)} = \sum_{j=1}^N \gamma_{j,k}^{(i)}$$

$$\left\{ \sum_{i=1}^N \sum_{k=1}^K \gamma_{i,k}^{(i)} \right\} N$$

$$A_{lim} = \sum_{i=1}^N \sum_{k=1}^K \gamma_{i,k}^{(i)}$$

$$\sum_{i=1}^N \left(\sum_{k=1}^K \sum_{m=1}^M \gamma_{i,k,m}^{(i)} \right)$$

$\eta_k \rightarrow$ soft count
ML
 $\downarrow$
 $(\hat{\mu}_k, \hat{\Sigma}_k)$
for Gaussian

e.g. Gaussians, similar to GMM $\rightarrow$ "weighted empirical mean" with weights $\tau_{i,k}^{(i)}$

$$\hat{\mu}_k = \frac{\sum_{i=1}^N \sum_{k=1}^K \gamma_{i,k}^{(i)} x_i^{(i)}}{\sum_{i=1}^N \sum_{k=1}^K \gamma_{i,k}^{(i)}}$$

Viterbi to compute $\arg \max_{z_{1:T}} p(z_{1:T} | \tilde{x}_{1:T})$
(max product)

Information Theory

KL divergence : for discrete dist. $p \neq q$

$$KL(p \parallel q) \triangleq \mathbb{E}_p \left[\log \frac{p(x)}{q(x)} \right] = \sum_{x \in \Omega} p(x) \log \frac{p(x)}{q(x)}$$

$$0 \cdot \log 0 = 0$$

$$\left(\lim_{x \rightarrow 0^+} x \log x = 0 \right)$$

[if $\exists x$ s.t. $q(x) = 0$
but $p(x) \neq 0$

$$\left[-p(x) \log q(x) = +\infty \right]$$

not zero $-\infty$

$$KL = \infty \text{ if } \text{supp}(p) \not\subseteq \text{supp}(q)$$

motivation from density estimation

recall statistical decision theory

(statistical) loss

would

$$L(p_\theta, \hat{q})$$

here, estimation of dist., say $\hat{q}$

$$\text{standard (MLE) loss is log-loss } L(p_\theta, \hat{q}) = \mathbb{E}_{x \sim p_\theta} \left[-\log \hat{q}(x) \right]$$

if use $\hat{q} = p_\theta$, then get

$$\sum_{x \in \Omega_x} -p_\theta(x) \log p_\theta(x) \triangleq H(p_\theta)$$

entropy of p_θ

is the q-p, then you

$$\sum_{x \in \Omega} p(x) \log p(x) = H(p)$$

entropy of p

excess loss for action $a = \hat{q}$

$$L(p, \hat{q}) - \min_{\hat{q}} L(p, \hat{q}) = \sum_{x \in \Omega} -p(x) \log \frac{\hat{q}(x)}{p(x)} = KL(p, \hat{q})$$

$\underbrace{\min_{\hat{q}} L(p, \hat{q})}_{L(p, p)}$

coding theory:

use length of code $\ell \sim -\log p(x)$

$\lg \triangleq \log_2 \rightarrow$ "bits"
 $\log_e \rightarrow$ "nats"

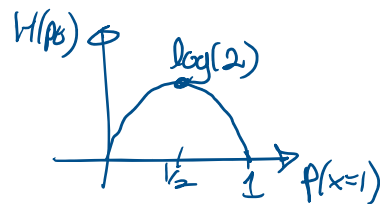
expected length of code: $\sum_x p(x) (-\log p(x))$ (entropy measured in bits)

KL divergence $\rightarrow$ interpreted as excess length cost (in terms of length of code) to use dist q to design code vs optimal dist (mc p)

example:

entropy of a Bernoulli

$$-p \log p - (1-p) \log (1-p)$$



entropy of a uniform dist. on k states

$$\sum_{x=1}^k -\frac{1}{k} \log \left(\frac{1}{k}\right) = \log k$$

(max entropy dist. over k states)

properties of KL

• $KL(p \parallel q) \geq 0$ $\leftarrow$ to show this, use Jensen's ineq. $f(\mathbb{E}X) \leq \mathbb{E}f(X)$ when f is convex

• KL is strictly convex in each of its arguments

i.e. $KL(p \parallel \cdot) : \Delta_K \subseteq \mathbb{R}^K \rightarrow \mathbb{R}$
 $\downarrow$ strictly convex

$$KL(\cdot \parallel q)$$

• not symmetric $KL(p \parallel q) \neq KL(q \parallel p)$
in general

$$KL(p \parallel p) = 0 \quad \forall p \in \Delta_K$$

commutational $KL(p \parallel q) + KL(q \parallel p)$

—

" in general "

symmetrized
version

$$\frac{1}{2}(KL(p||q) + KL(q||p))$$

" " "

" " "