

- today:
- MCMC
 - review Markov chains
 - M-H. alg.

MCMC - Markov chain Monte-Carlo

idea: is to relax indep assumption samples

to allow adaptive proposal dist.

ie. we'll run a chain $X_t | X_{t-1}$ s.t. $X_t \xrightarrow{t \rightarrow \infty} \text{in dist.}$
 $\downarrow$
 target dist. p

"stationary dist. of a chain"

then we can approximate

$$\mathbb{E}_p[f(x)] \text{ as } \frac{1}{T-T_0} \sum_{t=T_0+1}^T f(X_t)$$

T_0 is called "burn in period" $\rightarrow$ depends on "mixing time" of Markov chain

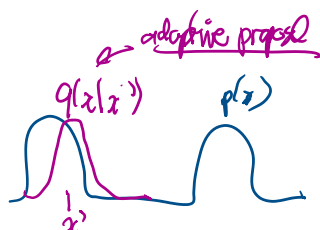
⊗ no need to thin the samples [ie. use Δt between samples to get more independence]

as this yields higher variance

$\rightarrow$ better to use all samples after T_0 to estimate μ

[unless it is too expensive]

motivation



before: samples $X^{(i)}$ i.i.d. q

$$\text{MCMC } X^{(t)} | X^{(t-1)} \sim q(\cdot | X^{(t-1)})$$

$\nearrow$
Markov transition probs

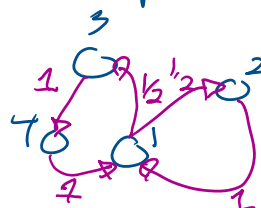
Review of (finite state space) Markov chain [$|X|=k$]

- as PGM



- there is also a transition prob. of view: use one node per state (probabilistic FSA)

e.g. $k=4$



[homogeneous M.C.]

↳ i.e. $P\{X_t = i \mid X_{t-1} = j\} = A_{ij}$ (no time dep.)

A is a $k \times k$ matrix s.t. $\mathbf{1}_k^T A = \mathbf{1}_k^T$
 "left-stochastic matrix" vector of ones of size k

⊛ (as in HMM) suppose $P\{X_{t-1} = j\} = (\pi)_j$

$$P\{X_t = i\} = \sum_j \underbrace{P\{X_t = i \mid X_{t-1} = j\}}_{A_{ij}} \underbrace{P\{X_{t-1} = j\}}_{\pi_j}$$

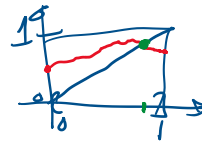
$$\begin{aligned} \pi_t &= A \pi_{t-1} \\ \Rightarrow \pi_t &= A^t \pi_0 \end{aligned}$$

stationary dist π of A is a dist. π s.t. $A\pi = \pi$

[note that this π is a right e-vector of A with e-value 1]

fact: every stochastic matrix has at least 1 stat. dist.

(by Brouwer's fixed pt. thm.)



def: irreducible Markov chain $\Leftrightarrow$ there exists a positive prob "path" from any i to any j (states)
 $\forall (i,j), \exists \text{ an integer } m_{ij} \text{ s.t. } (A^{m_{ij}})_{ji} > 0$

(by Perron-Frobenius thm. $\Rightarrow$ irreducible M.C., has a unique stat. dist.)
 [multiplicity of e-value 1 is 1]

⊛ in order to converge to it, we need aperiodicity as well

irreducible and aperiodic M.C. $\Leftrightarrow \exists$ an integer m s.t. $(A^m)_{ij} > 0$

aka regular M.C. (finite state space)

$$(A^m)_{ij} > 0 \quad \forall i,j$$

or ergodic M.C.

⊛ [note: a sufficient condition for an irreducible M.C. to]

the aperiodic is $\exists i$ s.t. $A_{ii} > 0$

example of a regular M.C. on k states ($k \geq 3$)

$$A = \begin{pmatrix} 0 & \frac{1}{2} & \frac{1}{2} \\ \frac{1}{2} & 0 & \frac{1}{2} \\ \frac{1}{2} & \frac{1}{2} & 0 \end{pmatrix} = \frac{1}{k-1} (11^T - I)$$

$$A^2 = \frac{1}{(k-1)^2} (111^T - 211^T + I) = \frac{1}{(k-1)^2} [(k-2)11^T + I] \quad \text{for } k \geq 3 \quad \boxed{\text{this is } > 0 \text{ for } m=2}$$

[but, for $k=2$ it is not aperiodic $A^2 = I$ $A = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$]

thm. if a finite M.C. is ergodic (regular)

then $\exists$ a unique stat. dist. π

and for any starting dist π_0 , $\lim_{t \rightarrow \infty} A^t \pi_0 = \pi$

the speed of convergence is related to the mixing time τ of the chain

$$\tau \triangleq \frac{1}{1 - |\lambda_2(A)|}$$

2nd largest e-value of A

$$\|A^t \pi_0 - \pi\|_1 \leq C \exp(-t/\tau)$$

after τ steps, error decreases by $\frac{1}{2}$

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⊗ intuition (from linear algebra) [informal argument]

simplify case, suppose A is symmetric & psd

by spectral thm, A is diagonalizable with orthogonal U

$$A = U \Sigma U^T \quad \text{with} \quad \Sigma = \begin{pmatrix} \lambda_1 & & 0 \\ & \ddots & \\ 0 & & \lambda_k \end{pmatrix}$$

$U \rightarrow$ orthonormal basis of e-vectors of A

$$U U^T = U^T U = I \quad U = (u_1 \dots u_k)$$

let α_0 be coordinates of π_0 in U basis

[by Perron-Frobenius thm
 $\lambda_1 = 1 > |\lambda_2| \geq \dots \geq \lambda_k$
 $\downarrow$
 $u_1 = \frac{\pi}{\|\pi\|_2} \quad A\pi = \pi$

let α_0 be coordinates of π_0 in U basis

i.e. $\pi_0 = U\alpha_0$ ($\alpha_0 = U^T\pi_0$)

$$A^t \pi_0 = (U \underbrace{\Sigma^t U^T}_{I}) \underbrace{U^T \pi_0}_{\alpha_0}$$

$$= U \Sigma^t \alpha_0$$

$$\Sigma^t = \begin{pmatrix} 1 & & \\ & \lambda_1^t & \\ & & \ddots & \\ & & & \lambda_k^t & \\ & & & & 0 \end{pmatrix}$$

$$= (\alpha_0)_1 1^t u_1 + (\alpha_0)_2 \lambda_2^t u_2 + \dots + (\alpha_0)_k \lambda_k^t u_k$$

$$\|A^t \pi - \pi\|_1 = \frac{(\alpha_0)_1}{\|u_1\|_1} = 1 \quad \text{not true in general}$$

$$\|(\alpha_0)_2 \lambda_2^t u_2 + \dots\| \leq C |\lambda_2|^t$$

exist e-gap
 $\lambda_2 = 1 - \epsilon_1$

$$\epsilon_1 \triangleq 1 - |\lambda_2|$$

$$|x| \leq \exp(x) \quad \forall x$$

$$|\lambda_2| \leq \exp(-\epsilon_1)$$

$$|\lambda_2|^t \leq \exp(-t\epsilon_1)$$

$$\frac{1}{t}$$

$$\Rightarrow \uparrow = \frac{1}{t |\lambda_2|}$$

⊛ mixing time is often (usually) exponentially large ⊛

⊛ How do we design A s.t. $A^t \pi_0 \rightarrow \pi$?

one "easy way"

reversible M.C.

$\Leftrightarrow$

$\exists$ dist. π s.t.

$$A_{ij} \pi_j = A_{ji} \pi_i \quad \forall (i,j)$$

"detailed balance equation"

this is sufficient condition

It means $\mathbb{P}\{X_t = i\} = \pi_i$

to get $A\pi = \pi$

detailed balance

then $\mathbb{P}\{X_t = i, X_{t-1} = j\}$

$$\text{proof: } (A\pi)_i = \sum_j A_{ij} \pi_j = \sum_j A_{ji} \pi_i = \pi_i \left(\sum_j A_{ji} \right) = \pi_i \cdot 1 //$$

Metropolis-Hastings alg:

goal $\rightarrow$ construct a M.C. with stat dist. $p(x)$ [can target]

[assume $p(x) > 0 \quad \forall x$]

use some proposal $q(x'|x)$

accept new state x' with prob. $\alpha(x'|x)$

$$\alpha(x'|x) \triangleq \min \left\{ 1, \frac{q(x|x') p(x')}{q(x'|x) p(x)} \right\}$$

no dependence on normalization of p

accept new state x' with prob. $\alpha(x'|x) \triangleq \min\left\{1, \frac{q(x|x')p(x')}{q(x'|x)p(x)}\right\}$
 if reject $\rightarrow$ stay in same state x
 [this is still new sample]
 vs.
 rejection sampling where only "accepted states" are samples
 acceptance ratio is satisfy detailed balance

M.H. alg:

start at $x^{(0)}$

important design choice

for $t=1, \dots$
 • propose $x^{(t)} \sim q(x'|x^{(t-1)})$
 flip a coin, with prob. $\alpha(x^{(t)}|x^{(t-1)})$ to be 1
 • if accept (coin=1)
 let $x^{(t)} = x^{(t)}$
 or
 let $x^{(t)} = x^{(t-1)}$
 end for

note: for symmetric proposal $q(x'|x) = q(x|x')$; always accept if $p(x') \geq p(x)$
 [Metropolis alg.] $\leadsto$ like noisy hill-climbing

[verify as exercise that M.H. satisfies detailed balance eq. with $\pi = p$ ^{target}]

⊛ for convergence: if M.H. chain is ergodic, then converge to unique stat. dist. p

sufficient conditions

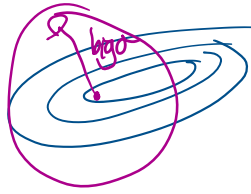
for irreducible $q(x'|x) > 0 \forall x' \neq x \in X$
 for aperiodicity $\begin{cases} \text{either } q(x|x) > 0 \text{ for some } x \in X \\ \text{or } \alpha(x'|x) < 1 \text{ for some } x \in X \text{ and } x' \text{ s.t. } q(x'|x) > 0 \end{cases}$
 $\downarrow$
 $A_{ii} > 0$ for some i

⊛ aside: it is still ok change proposal with time

$q_t(x'|x)$ as long as choice of q_t does not depend on $x^{(t+1)}$
 then convergence theory above will go through

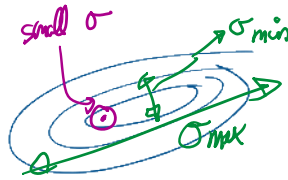
slow mixing example

suppose p is a $N(\mu, \Sigma)$



→ high prob rejection

$$q(x'|x) = N(x'|x, \sigma^2 I)$$



'small moves' $\Rightarrow$ many steps needed

here best mixing time related to $\frac{\sigma_{\max}}{\sigma_{\min}}$

reference for mixing times:

Markov Chains and Mixing Times
David A. Levin, Yuval Peres, Elizabeth L. Wilmer
<https://pages.uoregon.edu/dlevin/MARKOV/markovmixing.pdf>