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The Repair vs Replacement Problem : A Stochastic Control Approach

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SUMMARY

This paper is concerned with a repair vs replacement problem for a system observed in continuous time and subject to random failure. When the system fails, a decision is taken, based on the observed values of the (random) repair cost and system's age, to either perform a minimal repair or to replace the system by a new one. At any time, the system may also be replaced preventively, and the aim of the decision maker is to minimize the total expected discounted cost. A dynamic programming approach is proposed to solve this problem. Sufficient conditions under which a repair-limit rule is optimal are given, an efficient algorithm is devised for the computation of such an optimal policy, and a numerical illustration is worked out.

KEY WORDS: Dynamic programming Stochastic control Optimal replacement Repair-limit

RESUME

Cet article présente un modèle d'entretien d'équipement où à chaque intervention on peut soit remplacer le système soit effectuer une réparation minimale. Une approche basée sur des développements récents de la programmation dynamique stochastique permet de caractériser la solution et fournit un algorithme efficace pour calculer une stratégie optimale.

1. Introduction

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The aim of this paper is to propose a very general formulation of the so-called repair vs replacement problem for a system observed in continuous time, to develop a dynamic programming approach for its solution, to find conditions under which a repair-limit rule is optimal and, finally, to devise an efficient algorithm for the computation of such an optimal policy.

The repair vs replacement problem is a variant of the classical preventive replacement problem for an age-dependent failure-prone system^{1,2,10,16} where, at failure time, there is a possibility either to perform a minimal repair or to replace the system by a new one. The model presented in sections 2 and 3 generalizes earlier studies^{3,7,9} along several lines:
(i) a very general cost structure is assumed; (ii) no a priori assumptions on the form of the optimal policy are made; (iii) the problem is formulated in the general framework of discrete-event dynamic programming with continous state variables. In section 4, we explore the conditions under which a repair-limit rule is an optimal policy. In a repair-limit rule, a decision to replace the system is taken either at a predetermined age limit, or at failure time provided that the cost of a minimal repair is higher than a predetermined limit which is generally dependent on the age of the system. In section 5, we propose an efficient algorithm to compute an ϵ -optimal policy, for the case where a repair-limit rule is optimal. In section 6, the algorithm is illustrated on a numerical example. The appendix contains the mathematical developments involved in establishing the results of the paper.

This application is based on recent theoretical developments in dynamic programming for discrete-event systems with state variables in Borel spaces^{5,13,14,15}, and on an advanced implementation of a successive approximations algorithm in such a general framework¹². This approach has already permitted a complete solution of a continuous-time group preventive replacement problem¹², and the present paper gives another example of its usefulness for the treatment of complex maintenance problems. In this general modern framework, discrete event stochastic systems can be modelled without an arbitrary a priori time and state discretization, as it was often the case in previous stochastic dynamic programming implementations for maintenance applications^{4,10,12}.

2. Modeling the repair vs replacement problem

Consider a failure-prone system with an operating cost rate $c_o(s)$ and a failure rate $\lambda(s)$ at age s. If the system fails at age s, a failure cost $c_f(s)$ is incurred, and the system is either replaced at a cost $c_r(s)$, or repaired at a cost r, which is the sample value of a random variable whose probability law is defined by the conditional probability measure $R(\cdot \mid s)$. We assume that the decision to replace or repair is taken after having observed the actual value of r. A (minimal) repair brings the system back to an operational state without changing its age s, while a replacement brings the age of the system back to 0. At any age s, one may replace preventively the system, at a cost $c_r(s)$, in order to avoid possible failure costs. The system operates on an infinite time horizon and the costs are discounted at a rate $\rho > 0$.

The state of the system at time t is, by convention, its age at time t^+ ; i.e. at any intervention time (replacement or repair) the state s represents the age of the system immediately after the intervention. An intervention is triggered either by a failure or by a planned preventive replacement. The state space is thus $S = [0, \infty)$. A controller action can be represented by a pair $a = (d, \gamma)$, where $d \in [0, \infty)$ is the time delay till the next planned preventive replacement, and $\gamma : S \times [0, \infty) \to \{0, 1\}$ is a (measurable) binary function, called the repair vs replacement rule, such that if a failure occurs at age s and the observed repair cost is r, a replacement is decided if $\gamma(s,r) = 1$, and a repair is decided if $\gamma(s,r) = 0$. Let A denote the set of all possible controller actions.

The evolution of the system is then described as follows. Call stage n the occurence of the n-th intervention. At initial stage 0, the system is in state 0. Entering stage n at time τ_n , $n \geq 0$, the controller observes the age s_n of the system and then chooses an action $a_n = (d_n, \gamma_n)$. Given the failure process, the state s_n and the chosen action, a random transition defines (i) the next intervention time τ_{n+1} , which corresponds to stage n+1 and is the minimum between the time $\tau_n + \zeta$ of the next failure and the time $\tau_n + d_n$ of the next planned preventive replacement; and (ii) the state s_{n+1} of the system at stage n+1, which is zero if the system is replaced or $s_n + \zeta$ if the system is repaired.

This system is thus modelled as a Markov Renewal Decision Process^{8,13,14,15} with continuous state and action sets. In order to complete the definition of this process, we introduce the survival function

$$\overline{F}(\zeta) = e^{-\int_0^{\zeta} \lambda(t)dt},\tag{1}$$

the conditional survival function given state s,

$$\overline{F}(s+\varsigma\mid s) = \frac{\overline{F}(s+\varsigma)}{\overline{F}(s)} \tag{2}$$

and the conditional probability density function of the age of failure

$$f(s+\varsigma\mid s)=-\frac{\partial}{\partial\varsigma}\overline{F}(s+\varsigma\mid s)=\lambda(s+\varsigma)\overline{F}(s+\varsigma\mid s), \qquad f(\varsigma)=f(\varsigma\mid 0). \tag{3}$$

One can then define the expected one-stage transition cost given state s and action $a = (d, \gamma)$ as

$$g(s,d,\gamma) = \int_0^d \left[\int_0^{\varsigma} e^{-\rho t} c_o(s+t) dt + e^{-\rho \varsigma} \left(c_f(s+\varsigma) + \int_0^{\infty} w(s+\varsigma,r) R(dr \mid s+\varsigma) \right) \right] f(s+\varsigma \mid s) d\varsigma + \left[\int_0^d e^{-\rho t} c_o(s+t) dt + e^{-\rho d} c_r(s+d) \right] \overline{F}(s+d \mid s)$$

$$(4)$$

where

$$w(s+\zeta,r) = \begin{cases} c_r(s+\zeta) & \text{if } \gamma(s+\zeta,r) = 1, \\ r & \text{otherwise.} \end{cases}$$
 (5)

A (stationary) policy is a map $\mu: S \to A$ which associates with a state s an action $a = \mu(s)$. For the ease of exposition, we omit all mathematical details relating to the measurability questions and refer the reader to the appendix, where these questions are addressed. The term admissible will always refer to the minimum measurability properties needed for the existence of the integrals (including those defining the mathematical expectations). An admissible policy μ generates a Markov renewal stochastic process with value in $S \times A$, with an expected total discounted cost, given initial state $s_0 = s$:

$$V_{\mu}(s) = E_{\mu,s} \left[\sum_{n=0}^{\infty} e^{-\rho \tau_n} g(s_n, a_n) \right]. \tag{6}$$

The problem is to find an optimal policy μ^* , such that

$$V_{\mu^*}(s) = V_*(s) \stackrel{\text{def}}{=} \inf_{\mu} V_{\mu}(s)$$
 (7)

or an ϵ -optimal policy μ , such that for a given $\epsilon > 0$,

$$V_{\mu}(s) \le V_{*}(s) + \epsilon \tag{8}$$

for all s in S.

Remark: An equivalent model could be defined with the pair (s, r) taken as the state variable, s being the age of the system at t^- , i.e. just before the intervention, and r being the observed repair cost at failure time. This would simplify the action space, but increase the dimension of the state set (from one-dimensional to two-dimensional). From the computational viewpoint, it is much preferable to keep the smallest possible state set. Our choice also makes easier the theoretical developments of the following sections.

3. The dynamic programming formalism

When the state set is countable, it is well known that under some regularity conditions, a dynamic programming approach can be used for the solution of an infinite horizon discounted Markov renewal decision process (or semi-Markov decision process)^{8,10}. For the problem defined in section 2, one has two additional difficulties, namely a continuous state space and absence of the classical regularity conditions^{5,8}, which ask for the expected time between any two successive transitions to be bounded away from zero. A dynamic programming formalism adapted to this situation has been developed recently^{5,13,14,15}, and is mathematically more intricate because of delicate measurability problems. In this section, we use a dynamic programming approach in operator form and provide conditions under which a successive application of the operator converges to its fixed point.

For a given action $a=(d,\gamma)$, define the operator $H(V)(\cdot,d,\gamma)$ which transforms any admissible cost-to-go function $V:S\to [0,\infty)$ into another admissible cost-to-go function :

$$H(V)(s,d,\gamma) = g(s,d,\gamma)$$

$$+ \int_0^d \int_0^\infty V\left((1-\gamma(s+\varsigma,r))(s+\varsigma)\right) R(dr\mid s+\varsigma)e^{-\rho\varsigma}f(s+\varsigma\mid s)d\varsigma$$

$$+ V(0)e^{-\rho d}\overline{F}(s+d\mid s). \tag{9}$$

With any admissible policy μ , we associate the operator T_{μ} defined by

$$T_{\mu}(V)(s) = H(V)(s, \mu(s)).$$
 (10)

Finally, we define the dynamic programming operator T by

$$T(V)(s) = \inf_{\mu} T_{\mu}(V)(s).$$
 (11)

Value iteration (also called successive approximations) consists in applying the operator T repeatedly on an approximation of the true cost-to-go function. Starting from the zero function, the value iteration algorithm converges to V_* under adequate regularity conditions. We will see in *Proposition 1* that *Condition 1* below is sufficient for $T^n(V_0)$ to converge uniformly to V_* as n tends to infinity, where V_0 denotes the zero function and T^n is the n-fold composition of the operator T.

CONDITION 1.

(i) There is a constant L > 0 such that for all s in S,

$$\min\left(c_r(s), \quad c_o(s)/\lambda(s) + c_f(s) + \int_0^\infty \min(r, c_r(s)) R(dr \mid s)\right) \geq L; \tag{12}$$

(ii) There exists an admissible policy μ_1 and two constants $U \ge 0$ and $0 \le b < 1$ such that for all $s \ge 0$,

$$g(s, \mu_1(s)) \leq U \text{ and } \alpha(s, \mu_1(s)) \leq b, \tag{13}$$

where

$$\alpha(s,d,\gamma) = \int_0^d e^{-\rho\varsigma} f(s+\varsigma \mid s) d\varsigma + e^{-\rho d} \overline{F}(s+d \mid s)$$
 (14)

denotes the one-stage expected discount factor when one chooses action (d, γ) in state s.

By a staightforward calculation (see the appendix), the one-stage cost can be rewritten differently as

$$g(s,d,\gamma) = \int_0^d \left[c_o(s+\varsigma) + \lambda(s+\varsigma) \left(c_f(s+\varsigma) + \int_0^\infty w(s+\varsigma,r) R(dr \mid s+\varsigma) \right) \right]$$

$$e^{-\rho\varsigma} \overline{F}(s+\varsigma \mid s) d\varsigma + c_r(s+d) e^{-\rho d} \overline{F}(s+d \mid s), \tag{15}$$

which may facilitate the interpretation of condition 1. Equation (13) states that for some policy μ_1 , the one-stage cost is bounded and the expected duration between any two stages is bounded away from zero. Equation (12) implies that the one-stage expected cost divided by the one-stage expected discount factor is bounded away from zero. More specifically, for each (s, d, γ) in $S \times A$, we have

$$g(s,d,\gamma) \geq \alpha(s,d,\gamma)L.$$
 (16)

This shows that condition 1 also implies that this model is a so-called version LC of the Markov renewal decision process model introduced in references 13, 14, 15, for which the following result holds true:

PROPOSITION 1.

Under condition 1, we have:

- (i) T(V) = V iff $V = V_*$;
- (ii) $\lim_{n\to\infty}\sup_{s\in S}|T^n(V)(s)-V_*(s)|=0$ for all admissible V;
- (iii) A policy μ is optimal iff $T_{\mu}(V_*) = V_*$.

Proof: See references 14 and 15. ■

The repair vs replacement problem has thus been modelled in such a way that it lends itself to a value iteration dynamic programming approach for its solution. Clearly this modeling approach considers an a priori very large action space. In the next section, it is shown that, for an interesting class of systems, the policy space can be reduced to a much smaller class of decision rules, namely the *repair-limit* rules. For these forthcoming developments, it will be convenient to define for each d in $[0, \infty)$ the operator $J(V)(\cdot, d)$ by

$$J(V)(s,d) = H(V)(s,d,\gamma_V) = \inf_{\gamma} H(V)(s,d,\gamma)$$
 (17)

where

$$\gamma_{V}(s+\zeta,r) = \begin{cases} 1 & \text{if } r+V(s+\zeta) \geq c_{r}(s+\zeta)+V(0); \\ 0 & \text{otherwise.} \end{cases}$$
 (18)

One thus has

$$T(V)(s) = \inf_{0 \le d \le \infty} J(V)(s, d). \tag{19}$$

4. Optimality conditions for repair-limit policies

A repair-limit policy is a rule under which the decision maker replaces the component either at a predetermined age d_* , or when it fails at age s and the repair cost is greater than a predetermined age-dependent limit $L_*(s)$; any other failure being corrected by a minimal repair. More formally, a repair-limit policy is a policy μ such that

$$\mu(s) = \begin{cases} (d_* - s, \gamma_*) & \text{if } s < d_*; \\ (0, \gamma_*) & \text{otherwise} \end{cases}$$
 (20)

where

$$\gamma_*(s,r) = \begin{cases} 1 & \text{if } r > L_*(s) \text{ or } s \ge d_*; \\ 0 & \text{otherwise;} \end{cases}$$
 (21)

and where d_* is a fixed non-negative real number and $L_*: [0, d_*] \to [0, \infty)$ is a given measurable function. This definition generalizes the one introduced by Drinkwater and Hastings⁹. In this section, we characterize a class of *repair vs replacement* problems for which an optimal policy is of the *repair-limit* type.

CONDITION 2.

Assume that $c_r(\cdot)$ and $\lambda(\cdot)$ are nondecreasing functions of the age s, that the derivative $c_r'(\cdot)$ of $c_r(\cdot)$ exists and that, for any nondecreasing function $q:[0,\infty)\to[0,\infty)$, bounded above by $c_r(\cdot)$, the following functionals

$$\Lambda_1(s,q(s)) = c_0(s) + \lambda(s) \left[c_f(s) + \int_0^{c_r(s) - q(s)} (q(s) + r - c_r(s)) R(dr \mid s) \right]$$
(22)

$$\Lambda_2(s, q(s)) = \Lambda_1(s, q(s)) + c_r'(s) - \rho c_r(s)$$
(23)

are nondecreasing in s.

We claim that the systems for which condition 2 holds admit optimal policies which are *repair-limit*. More precisely, if one calls B_+ the class of nondecreasing admissible cost-to-go functions $V(\cdot)$, and if one defines the operators $G_1: B_+ \to B_+$ and $G_2: B_+ \to B_+$ by

$$G_1(V)(s) = \Lambda_1(s, \min(c_r(s), V(s) - V(0)))$$
(24)

$$G_2(V)(s) = \Lambda_2(s, \min(c_r(s), V(s) - V(0))) - \rho V(0)$$
(25)

and the functional

$$l(V) = \inf \{ s \ge 0 \mid G_2(V)(s) \ge 0 \}$$
 (26)

then the following result, whose proof is given in the appendix, holds.

PROPOSITION 2.

Under condition 2, one has:

(i) for each V in B_+ , T(V) is in B_+ and

$$T(V)(s) = \begin{cases} J(V)(s, l(V) - s) & \text{if } s < l(V); \\ J(V)(s, 0) = c_r(s) + V(0) & \text{if } s \ge l(V); \end{cases}$$
(27)

(ii) V_* is in B_+ and the policy μ_* defined by

$$\mu_*(s) = \begin{cases} (d_* - s, \gamma_*) & \text{if } s < d_*; \\ (0, \gamma_*) & \text{otherwise;} \end{cases}$$
 (28)

where $d_* = l(V_*)$ and

$$\gamma_*(s,r) = \begin{cases} 1 & \text{if } r \ge L_*(s) \stackrel{\text{def}}{=} V_*(0) + c_r(s) - V_*(s); \\ 0 & \text{otherwise;} \end{cases}$$
(29)

is optimal.

This result expresses the optimal maintenance policy as a repair-limit rule. Of course both the age limit d_* and the repair limit function $L_*(s)$ depend on the optimal cost-to-go function $V_*(\cdot)$. Therefore, the above theorem does not specify completely the optimal policy, it only gives an information on the form of the optimal policy. This information will be exploited in the algorithm proposed in section 4.

Before leaving this section we can show that the class of systems which satisfy *condition* 2 encompasses the Drinkwater and Hastings model⁹ as well as the classical age replacement model².

EXAMPLE 1.

For the model considered in reference 9, one has $c_f = c_o = 0$, c_r is a constant, $\lambda(s)$ is nondecreasing in s, and $R(\cdot \mid s)$ is an exponential distribution with mean m(s), where m(s) is nondecreasing in s. Eq. (22) then becomes

$$\Lambda_1(s, q(s)) = \lambda(s) \left[q(s) - c_r + m(s) \left(1 - e^{-(q(s) - c_r)/m(s)} \right) \right]$$
(30)

and condition 2 is easily seen to hold.

EXAMPLE 2.

Consider a model where minimal repair is not allowed, or where the repair cost is always (with probability 1) larger than the replacement cost, that is $R([0, c_r(s)] \mid s) = 0$ for all s. The integral in equation (22) then vanishes to 0 for all s and q, and one obtains

$$\Lambda_1(s,q(s))=c_o(s)+\lambda(s)c_f(s).$$

The optimal policy in that case, under condition 2, is age replacement^{2,16}.

5. Computing an optimal repair-limit policy

We now provide an efficient computational algorithm for the case where conditions 1 and 2 are satisfied. It is an adaptation of the more general algorithm proposed in reference 12, and it operates as follows.

ALGORITHM

- 1. Compute U, L and b satisfying condition 1, and h := U/(1-b);
- 2. Choose $\epsilon > 0$ and V in B_+ ;
- 3. Let $V_{-}(\cdot) := 0$, $V_{+}(\cdot) := h$;

repeat

- 4. Compute l(V) satisfying (26);
- 5. Compute T(V)(s) = J(V)(s, l(V) s) at a finite number of points in the interval [0, l(V)], including 0 and l(V);
- 6. Define $W(\cdot)$ in B_+ as an approximation to $T(V)(\cdot)$ on [0, l(V)] and equal to $c_r(\cdot) + V(0)$ on $(l(V), \infty)$;
- 7. Compute (or estimate) nonnegative real values δ^- , δ^+ , γ^- and γ^+ such that for all s,

$$-\delta^- \le T(V)(s) - W(s) \le \delta^+;$$

 $-\gamma^- \le W(s) - V(s) \le \gamma^+;$

8. Let $V_-(\cdot):=\max\left(V_-(\cdot),\,W(\cdot)-(\delta^-+\gamma^-)4h/L\right);$ $V_+(\cdot):=\min\,\left(V_+(\cdot),\,W(\cdot)+(\delta^++\gamma^+)4h/L\right);$

- 9. Let $V(\cdot) := W(\cdot)$;
- 10. Display any desired values;

until $V_{+}(s) - V_{-}(s) < \epsilon$ for all s;

- 11. Display $l(V)(\cdot)$, $W(\cdot)$ and any other desired statistics;
- 12. Stop.

PROPOSITION 3

After any iteration of the algorithm, one has $V_{-}(\cdot) \leq V_{+}(\cdot) \leq V_{+}(\cdot)$. Further, if the error in the approximation of T(V) converges to 0, that is if the successive values of δ^{-} and δ^{+} converge to 0, then $V_{+}(\cdot) - V_{-}(\cdot)$ converges to 0 uniformly and the algorithm terminates after a finite number of iterations.

Proof: It follows from Theorem 3.3 in reference 12. ■

Some remarks should be made concerning this algorithm. Firstly, the choice of the policy μ_1 is somewhat arbitrary, but in order to obtain tighter bounds for V_* , one should choose it to minimize h. Secondly, any method of interpolation or approximation may be used to obtain $W(\cdot)$ in step 6, and the method can also vary between iterations. Interpolation could be used if T(V)(s) is computed with good precision; otherwise approximation (e.g. least squares) is recommended. Thirdly, as explained in ref. 12, the bounds in step 7 can be exact and conservative, derived by using the monotonicity properties of $V(\cdot)$, $T(V)(\cdot)$ and $W(\cdot)$, or can be estimations. For instance, when T(V) is reasonably smooth (which is usually the case), one can perform step 7 by recomputing T(V) on a much finer regular mesh, computing the actual error of approximation at these new points, and taking the smallest and largest of these errors as estimates of δ^- and δ^+ respectively.

6. A numerical illustration

Let $\rho = 0.1$, $c_r(s) = 20$, $c_f(s) = 5$, $\lambda(s) = \min\{0.02s, 0.2\}$, $c_o(s) = \min\{s + e^{-s}, 10 + e^{-10}\}$ (a system of age greater than 10 is assumed to be equivalent to a system of age 10), and the distribution function of the repair cost be defined by $R([0,r] \mid s) = 1 - e^{r/2}$ for all $s \geq 0$ and $r \geq 0$. The failure rate $\lambda(s)$ corresponds (for $s \leq 10$) to a Weibull distribution, and the repair cost has an exponential distribution with mean 2. Consider the policy μ_1 defined by

$$\mu_1(s)=(d_1,\,\gamma_1)$$

where $d_1 = 10$ and $\gamma_1(\cdot) \equiv 1$. Eq.(13) is verified for b = 0.6833, U = 48.5, and one has h = U/(1-b) = 153.14, whereas eq.(12) is verified for L = 20. Condition 2 is easily verified, and the algorithm proposed in the previous section can be used.

A computer implementation has been done, in FORTRAN, and the program allows changing the evaluation points of T(V) or the method of approximation after any iteration. The initial $V(\cdot)$ is set to zero in step 2. In step 4, a sequential search is used to localize l(V) approximately, then binary search is used to find its exact position. In step 5, T(V) is computed at n points, evenly spaced on the interval [0, l(V)]. The integrals are evaluated using the Simpson's rule, with an integration step small enough to obtain a negligible quadrature error. $W(\cdot)$ is defined on [0, l(V)] by spline interpolation, using splines of arbitrary order, and where the knots are chosen according to equation (10) on p. 219 of reference 6.

After 30 iterations with n=10 evaluation points and splines of order k=2 (i.e. piecewise-linear interpolation), one obtains l(V)=6.7047 and V(0)=56.352. After 5 more iterations with k=2 and n=25, one obtains l(V)=6.7056; and 5 more iterations with k=4 (cubic splines) and n=25 yield l(V)=6.7058 and V(0)=56.365. Finally, after 10 more iterations with cubic splines and n=50, one still has the same values. T(V) can then be recomputed on a much finer grid (500 uniformly spaced points) to estimate δ^- and δ^+ , which yields the estimate bounds on V_* :

$$W(\cdot) - 0.00023 = V_{-}(\cdot) \le V_{*}(\cdot) \le V_{+}(\cdot) = W(\cdot) + 0.00009.$$

The relative error on V_* is thus negligible. The repair-limit at age s is $L_*(s) = c_r(s) - V_*(s) + V_*(0)$, and this function is displayed in *Figure 1*. Here, the repair-limit is decreasing in s, but notice that this is not true in general, when $c_r(\cdot)$ is not constant.

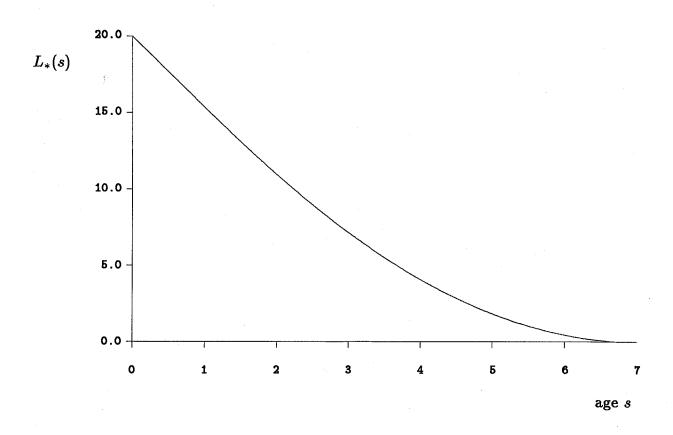


FIGURE 1. The optimal repair-limit function

Appendix

In this appendix, we state the measurability conditions sufficient for the dynamic programming formalism to be well defined and for *Proposition 1* to hold. We also derive eq.(15), and we prove *Proposition 2*.

A.1 Measurability conditions

Assume $c_f(\cdot)$, $c_r(\cdot)$, $c_o(\cdot)$ and $\lambda(\cdot)$ to be four Borel measurable functions from $[0, \infty)$ to $[0, \infty)$. R is assumed to be a continuous Borel measurable stochastic kernel^{5,13} on $[0, \infty)$ given $[0, \infty)$, i.e. a family $\{R(\cdot \mid s), s \geq 0\}$ of probability measures on $[0, \infty)$ such that for every Borel set $B \subseteq [0, \infty)$, $R(B \mid \cdot)$ is a continuous (and Borel measurable) function. The state space is $S = [0, \infty)$ and the action space is $A = [0, \infty] \times Z$ where Z is the class of Borel measurable binary functions $\gamma: S \times [0, \infty) \to \{0, 1\}$. An admissible policy is a universally measurable^{5,13} mapping $\mu: S \to A$. An admissible cost-to-go function is a bounded lower semi-analytic function $V: S \to [0, \infty)$. For any such function, the expressions in (9-11) are well defined^{5,13}.

Notice that more general kind of policies could have been considered, where μ is random and may depend on all past history. However, it could be shown^{5,14} that such an extension raises no additional gain on the optimal cost function.

A.2 Derivation of eq.(15)

Using Fubini's theorem, one has

$$\int_0^d \left(\int_0^{\varsigma} e^{-\rho t} c_o(s+t) dt \right) f(s+\varsigma) s) d\varsigma = \int_0^d \int_t^d e^{-\rho t} c_o(s+t) f(s+\varsigma \mid s) d\varsigma dt$$

$$= \int_0^d e^{-\rho t} c_o(s+t) \left[\overline{F}(s+t \mid s) - \overline{F}(s+d \mid s) \right] dt$$

and from (4), using the fact that $f(s+\zeta \mid s) = \lambda(s+\zeta)\overline{F}(s+\zeta \mid s)$, one obtains (15).

A.3 Proof of Proposition 2

Let $V \in B_+$. One has

$$\frac{\partial}{\partial d}J(V)(s,d) = e^{-\rho d}\overline{F}(s+d\mid s) \left[c_o(s+d) + \lambda(s+d) \left[c_f(s+d) + \int_0^\infty \min\{r+V(s+d), c_r(s+d)+V(0)\} R(dr\mid s+d) \right] + c_r'(s+d) - \left[c_r(s+d)+V(0) \right] \left[\lambda(s+d)+\rho \right] \right]$$

$$= e^{-\rho d}\overline{F}(s+d\mid s)G_2(V)(s). \tag{31}$$

From eq.(23) in condition 2, $G_2(V)(s)$ is non-decreasing in s. It is thus negative for s < l(V), non-negative for s > l(V), and one has

$$\frac{\partial}{\partial d}J(V)(s,d)$$
 $\begin{cases} < 0 & \text{if } d < l(V) - s; \\ \geq 0 & \text{if } d > l(V) - s. \end{cases}$

J(V)(s,d) is thus decreasing in d on the interval $[0, \max(0, l(V) - s)]$, non-decreasing thereafter, and the infimum in (19) is attained for $d = \max(0, l(V) - s)$. Notice that $d = \infty$ if $l(V) = \infty$.

It remains to show that T(V) is non-decreasing on (0, l(V)). Let $0 \le s_1 < s_2 < l(V)$ and $d_2 = l(V) - s_2$. One has

$$T(V)(s_{2}) = J(V)(s_{2}, d_{2})$$

$$= \int_{s_{2}}^{s_{2}+d_{2}} e^{-\rho(s-s_{2})} \overline{F}(s \mid s_{2}) \left[G_{1}(V)(s) + \lambda(s)(c_{r}(s) + V(0))\right] ds$$

$$+ e^{-\rho d_{2}} \overline{F}(s_{2} + d_{2} \mid s_{2}) \left[c_{r}(s_{2} + d_{2}) + V(0)\right]. \tag{32}$$

For each s in $[s_2, l(V)]$, let Q(s) be the value in $[s_1, s)$ satisfying

$$\int_{s_1}^{Q(s)} (\lambda(t) + \rho) dt = \int_{s_2}^{s} (\lambda(t) + \rho) dt.$$
 (33)

Such a value uniquely exists, since $\int_{s_1}^{s'} (\lambda(t) + \rho) dt$ is continuous and strictly increasing in s', and

$$\int_{s_1}^s (\lambda(t)+\rho)dt > \int_{s_2}^s (\lambda(t)+\rho)dt \geq 0 = \int_{s_1}^{s_1} (\lambda(t)+\rho)dt.$$

One then has

$$\overline{F}(s \mid s_{2})e^{-\rho(s-s_{2})} = e^{-\rho(s-s_{2}) - \int_{s_{2}}^{s} \lambda(t)dt}
= e^{-\int_{s_{2}}^{s} (\lambda(t) + \rho)dt} = e^{-\int_{s_{1}}^{Q(s)} (\lambda(t) + \rho)dt}
= e^{-\rho(Q(s) - s_{1}) - \int_{s_{1}}^{Q(s)} \lambda(t)dt}
= \overline{F}(Q(s) \mid s_{1})e^{-\rho(Q(s) - s_{1})}.$$
(34)

Deriving both sides of (33) with respect to s and using the fact that $\lambda(\cdot)$ is non-decreasing, one obtains

$$\frac{\partial Q(s)}{\partial s} = \frac{\lambda(s) + \rho}{\lambda(Q(s)) + \rho} \ge 1. \tag{35}$$

Let $d_1 = Q(l(V)) - s_1$. From (32, 34, 35), since $G_1(V)(\cdot)$, $\lambda(\cdot)$ and $c_r(\cdot)$ are non-decreasing on [0, l(V)], and with the change of variable t = Q(s), one obtains

$$T(V)(s_{2}) \geq \int_{s_{2}}^{s_{2}+d_{2}} e^{-\rho(Q(s)-s_{1})} \overline{F}(Q(s) \mid s_{1}) \left[G_{1}(V)(Q(s)) + \lambda(Q(s))(c_{r}(Q(s)) + V(0)) \right] \left[(\lambda(s)+\rho)/(\lambda(Q(s)) + \rho) \right] ds$$

$$+ e^{-\rho d_{1}} \overline{F}(s_{1}+d_{1} \mid s_{1}) \left[c_{r}(s_{1}+d_{1}) + V(0) \right]$$

$$= \int_{s_{1}}^{s_{1}+d_{1}} e^{-\rho(t-s_{1})} \overline{F}(t \mid s_{1}) \left[G_{1}(V)(t) + \lambda(t)(c_{r}(t) + V(0)) \right] dt$$

$$+ e^{-\rho d_{1}} \overline{F}(s_{1}+d_{1} \mid s_{1}) \left[c_{r}(s_{1}+d_{1}) + V(0) \right]$$

$$= J(V)(s_{1},d_{1})$$

$$\geq T(V)(s_{1}). \tag{36}$$

Thus, for each V in B_+ , T(V) is in B_+ , and since B_+ is closed, $V_*(\cdot) = \lim_{n \to \infty} T^n(V)(\cdot)$ is also in B_+ . Taking $d_* = l(V_*)$, one then has

$$J(V_*)(s, \max(0, d_* - s)) = \inf_{0 < d < \infty} J(V_*)(s, d)$$
(37)

for all s in S, and from Proposition 1 (iii), this completes the proof. \blacksquare

References

3.5

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